

An Improved Algorithm to Obtain Initial Basic Feasible Solution for the Transportation Problem

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Abstract: This proposed method constitutes alternate method for initial solution in transportation problem. This method provides the initial solution. A transportation matrix is solved by a difference between maximum and next maximum element for each row and difference between maximum and minimum element for each column. In which the maximum value is marked and allocation is given to the least element.. Salient features of this method depict lesser calculation time, easy applicability. Depiction with examples provides easy understanding of this method.

Keywords: Transportation problem, supply, Demand, Vogels method, initial solution

1. Introduction

Transportation problem is most referred and analyzed aspect where there is always efforts to match the demand-supply prerogative or to put them in balance. With each demand there needs to be sufficient supply to maintain balance and create optimal solution for transportation. The efficient way would be to provide a satisfactory logistical solution for the demand and supply with minimum cost incurred.

Considering the decision variable X_{ij} of model in transportation the i^{th} supply at source to j^{th} demand in destination.

The below listed methods are used to deduce feasible solutions with initial transportation problem:

- 1) North West Corner Method
- 2) Least Cost Method
- 3) Vogel's Approximation Method

The overtly popular method employed on solving this is from MODI and stepping stone method. The former being widely discussed and articulated. This new method can be applied instead of the above three methods to obtain an initial basic feasible solution.

The beginning of assignment problem dates back to 1941 by Hitchcockin, the subsequent evolution of the methods put forth Koopmans in 1949 and Dantzig in 1951. Though the Simplex method forms the basis of the problem solution, this would however be devoid of being implemented as it would not suit to the environment. With more research on its way in 1954, Charnes and Cooper introduced Stepping Stone Method. This proved to be more efficient, thereby driving the goal of being optimal. This one step further laid foundation towards our grit on further optimizing eventually giving birth to Heuristic method being formulated by Kirca and Stair through the VAM from Goyal;s version in the year 1958.

This article is formed in sections and flows as below:
 Mathematical representation of transportation problem is depicted in Section 2, followed by algorithm in Section 3,

numerical examples and conclusion follows in Section 4 and 5 respectively.

2. Mathematical form for transportation problem

The LP problem as given below

$$\text{Minimize } Z = \sum_{i=1}^m \sum_{j=1}^n C_{ij} X_{ij}$$

Subject to the constraints

$$\sum_{j=1}^n X_{ij} \leq S_i \text{ For all } i$$

$$\sum_{i=1}^m X_{ij} \geq d_j \text{ For all } j$$

$$X_{ij} \geq 0$$

Transportation problem is considered to be balanced if

$$\sum_{i=1}^m S_i = \sum_{j=1}^n d_j$$

3. Algorithm

Step 1

Construct the matrix of a transportation problem from given problem. In case if the problem is unbalanced we make it balanced.

Step 2

Find the difference between maximum and next maximum in each row which is called as row penalty and difference between maximum and minimum in each column as column penalty and write it in the side and bottom. From that select the maximum value. From the selected row / column we need to allocate the minimum of supply/demand in the minimum element of the row or column. Eliminate by deleting the columns or rows corresponding to where the supply or demand is satisfied.

Step 3

Thus obtained table is then discussed. The process is continued to the remaining table till $m + n - 1$ cells are allocated with satisfaction to its supply and demand.

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Step 4

If obtained condition in step 2 is contrary, that is if there is tie in maximum value select that value which has least element. If there is tie in the least element then allocate the least element which has minimum supply/demand.

Step 5

Repeating the steps 2 to step 4 until satisfaction of all the supply and demand is met.

Step 6

The total minimum cost is calculated by

$$\text{Total cost} = \sum_{i=1}^m \sum_{j=1}^n C_{ij} x_{ij}$$

4. Numerical Example

1. Obtain an initial basic feasible solution for the following problem.

Factories	Distribution				Supply
		D ₁	D ₂	D ₃	
	F ₁	6	4	1	50
	F ₂	3	8	7	40
	F ₃	4	4	2	60
	Req.	20	95	35	

Solution

Step 1

Find the difference between the maximum and next maximum element in row-wise and difference between maximum and minimum in column-wise. For that row/column allocate the minimum supply/demand for the least element in that row/column.

Factories		D ₁	D ₂	D ₃	Supply	Row Penalty
	F ₁			35	15	
	F ₂	6	4	1	50	(2)
	F ₃	3	8	7	40	(1)
	Req.	4	4	2	60	(0)
	Column Penalty	20	95	35		

we take the row penalty as F₁ (6-4=2) and F₂ (8-7=1) and F₃ (4-4=0) and column penalty as D₁ (6-3=3) and D₂ (8-4=4) and D₃ (7-1=6).

Here the difference is maximum in D₃. Hence allocate to the smallest element with minimum supply/demand. Now delete D₃

Step 2

Factories		D ₁	D ₂	Supply	Row Penalty
	F ₁	6	4	15	(2)
	F ₂	20	3	40-20=20	(5)*
	F ₃	4	4	60	(2)
	Req.	20	95		
	Column Penalty	(3)	(4)		

we take the row penalty as F₁ (18-15=3) and F₂ (16-13=3) and F₃ (17-12=5) and column penalty as D₁ (16-12=4) and D₂ (17-10=7) and D₄ (18-11=7)

Here the difference is maximum in F₂. Hence allocate to the smallest element with minimum supply/demand.

Now delete D₁

Step 3

Factories		D ₂	Supply
	F ₁	15	15
	F ₂	20	20
	F ₃	60	60
	Req.	95	
	Column Penalty	(4)	

Step 4

Factories		D ₁	D ₂	D ₃	Supply
	F ₁	6	15	35	50
	F ₂	20	20	7	40
	F ₃	4	60	2	60
	Req.	20	95	35	

$$\text{Minimum cost} = 15 \times 4 + 35 \times 1 + 20 \times 3 + 20 \times 8 + 60 \times 4$$

$$\text{Minimum cost} = 555$$

Initial Solution

Proposed Method	Vogels Method
555	555

2. Obtain an initial basic feasible solution for the following problem.

Factories		Distribution				Supply
		D ₁	D ₂	D ₃	D ₄	
	F ₁	15	10	17	18	20
	F ₂	16	13	12	13	60
	F ₃	12	17	20	11	70
	Req.	30	30	40	50	

Solution

Step1

Find the difference between the maximum and next maximum element in row-wise and difference between maximum and minimum in column-wise. For that row/column allocate the minimum supply/demand for the least element in that row/column.

		D ₁	D ₂	D ₃	D ₄	Supply	Row Penalty
Factories	F ₁	15	10	17	18	20	(1)
	F ₂	16	13	40 12	13	60	20 (3)
	F ₃	12	17	20	11	70	(3)
	Req.	30	30	40 10	50		
	Column Penalty	(4)	(7)	(8)*	(7)		

we take the row penalty as F₁(18-15=3) and F₂(16-13 =3) and F₃ (17-12=5) and column penalty as D₁(16-12=4) and D₂ (17-10=7) and D₄ (18-11=7)

Here the difference is maximum in D₂ and D₄. Select the column D₂ which has least element 10. Hence allocate to the smallest element with minimum supply/demand.

Now delete F₁.

Step 2

		D ₁	D ₂	D ₄	Supply	Row Penalty
Factories	F ₁	15	20 10	18	20	(3)
	F ₂	16	13	13	20	(3)
	F ₃	12	17	11	70	(5)
	Req.	30	30 10	50		
	Column Penalty	(4)	(7)	(7)		

Step 3

		D ₁	D ₂	D ₄	Supply	Row Penalty
Factories	F ₂	16	13	13	20	(3)
	F ₃	12	17	50 11	70 20	(5)*
	Req.	30	10	50		
	Column Penalty	(4)	(4)	(2)		

we take the row penalty as F₂(16-13=3) and F₃ (17-12=5) and column penalty as D₁(16-12=4) and D₂ (17-13=4) and D₄ (13-11=2).

Here the difference is maximum in F₃. Hence allocate to the smallest element with minimum supply/demand.

Now delete D₄

Step 4

		D ₁	D ₂	Supply	Row Penalty
Factories	F ₂	10 16	10 13	20	(3)
	F ₃	20 12	17	20	(5)*
	Req.	30	10		
	Column Penalty	(4)	(4)		

Step5

		D ₁	D ₂	D ₃	D ₄	Supply
Factories	F ₁	15	20 10	17	18	20
	F ₂	10 16	10 13	40 12	13	60
	F ₃	20 12	17	20	50 11	70
	Req.	30	30	40	50	

$$\text{Minimum cost} = 20 \times 10 + 10 \times 16 + 10 \times 13 + 40 \times 12 + 20 \times 12 + 50 \times 11$$

$$\text{Minimum cost} = 1760$$

Initial Solution

Proposed Method	Vogels Method
1760	1760

3. Obtain an initial basic feasible solution for the following problem.

		Distribution				
		D ₁	D ₂	D ₃	D ₄	Supply
Factories	F ₁	4	19	22	11	100
	F ₂	1	9	14	14	30
	F ₃	6	6	16	14	70
	Req.	40	20	60	80	

Solution

Step-1

Find the difference between the maximum and next maximum element in row-wise and difference between maximum and minimum in column-wise. For that row/column allocate the minimum supply/demand for the least element in that row/column.

Factories		D ₁	D ₂	D ₃	D ₄	Supply	Row Penalty
	F₁	4	19	22	11	100	(3)
	F₂	1	9	14	14	30	(0)
	F₃	6	20 6	16	14	70 50	(2)
	Req.	40	20	60	80		
Column Penalty		(5)	(13)*	(8)	(3)		

we take the row penalty as **F₁** (22-19=3) and **F₂** (14-14=0) and **F₃** (16-14=2) and column penalty as **D₁** (6-1=5) and **D₂** (19-6=13) and **D₃** (22-14=8) and **D₄** (14-11=3)

Here the difference is maximum in **D₂**. Hence allocate to the smallest element with minimum supply/demand.
 Now delete **D₂**

Step 2

Factories		D ₁	D ₃	D ₄	Supply	Row Penalty
	F₁	40 4	22	11	100 60	(10)*
	F₂	1	14	14	30	(0)
	F₃	6	16	14	50	(2)
	Req.	40	60	80		
Column Penalty		(5)	(8)	(3)		

we take the row penalty as **F₁** (22-11=10) and **F₂** (14-14=0) and **F₃** (16-14=2) and column penalty as **D₁** (6-1=5) and **D₃** (22-14=8) and **D₄** (14-11=3)

Here the difference is maximum in **F₁**. Hence allocate to the smallest element with minimum supply/demand.
 Now delete **D₁**

Step-3

Factories		D ₃	D ₄	Supply	Row Penalty
	F₁	22	60 11	60	(10)*
	F₂	14	14	30	(0)
	F₃	16	14	50	(2)
	Req.	60	80 20		
Column Penalty		(8)	(3)		

we take the row penalty as **F₁** (22-11=10) and **F₂** (14-14=0) and **F₃** (16-14=2) and column penalty as **D₃** (22-14=8) and **D₄** (14-11=3)

Here the difference is maximum in **F₁**. Hence allocate to the smallest element with minimum supply/demand.
 Now delete **F₁**

Step-4

Factories		D ₃	D ₄	Supply	Row Penalty
	F₂	30 14	14	30	(0)
	F₃	30 16	20 14	50	(2)*
	Req.	60	20		
Column Penalty		(2)	(0)		

Step-5

Factories		Distribution				
		D ₁	D ₂	D ₃	D ₄	Supply
	F₁	40 4	19	22	60 11	100
	F₂	1	9	30 14	14	30
	F₃	6	20 6	30 16	20 14	70
Req.		40	20	60	80	

Minimum cost = $40 \times 41 + 60 \times 11 + 30 \times 14 + 20 \times 6 + 30 \times 16 + 20 \times 14$

Minimum cost = 2120

Initial Solution

Proposed Method	Vogels Method
2120	2170

5. Conclusion

The above discussed method on initial solution gives another paradigm on solutions with more efficient way. This method further qualifies to be more efficient and initial in solving the problems being a game to further research and development through this. This proposed method gives result exactly or even lesser to VAM method. All necessary qualities of being time efficient, easy applicability etc., forms the core of being implemented successfully.

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