





area as it offers a new and complete framework to mechanize anomaly discovery and rule suppression in together centralized and distributed inheritance firewalls. Making of firewall policy relations is essential for analysis of the firewall policy and manipulative management techniques such as policy editing and anomaly discovery.

### 3.2 Proposed Architecture

The paper focuses on two key contributions. First, propose a privacy-preserving protocol for the detection of inter-firewall redundant rules in single firewall with respect to its adjacent firewall. Secondly, recognize all anomalies that might exist in a single-firewall or multi-firewall environment. In this there are different techniques for discovering these anomalies. Both functionalities are based on the rule-based segmentation technique.

#### 3.2.1 Conflict detection and resolution,

Conflict resolution techniques is used to enable a fine grained resolution of conflict with the assist of numerous effectual decision strategies regarding the risk evaluation of protected networks and the purpose of policy description.

#### 3.2.2 Redundancy discovery and removal.

This technique is used to detect redundancies in policies and removed to improve effective result.

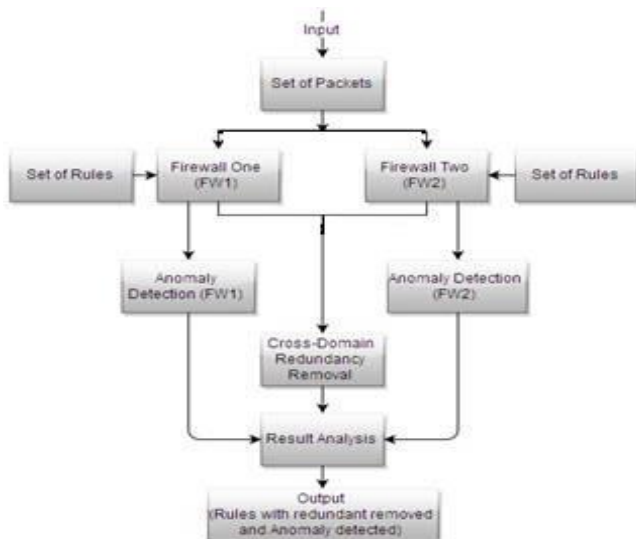


Figure 2: Proposed Architecture

### 3.3 Anomalies in Firewall Policies

Firewall policies are classified in to three different categories

#### Generalization-

If subset of packets is matched by the rule is also matched by subsequent rule but each rule taking different action then this rule is generalisation of one or set of previous rules.

#### Correlation-

If a one rule intersects with other rules and each rule defining a different actions then that rule correlated with other rules.

#### Redundancy-

In the same Firewall or same environment more than one rules are available and all those rules has the same effect then there is rule redundancy is defined.

### 3.4 Representation of Policy Anomaly

For effective anomaly detection, complete and accurate anomaly information diagnosis is represented in the Fig 3. When a more than one rule or number of rules interacts with each other, one overlapping relation may be associated with several rules. In anomaly detection technique there is possibility of one rule may overlap with multiple other rules and it and can be involved in other overlapping relations (overlapping segments).

Figure 3 shows a representation of policy anomalies. We can easily determine which rules are covered by a which segment, and which segments are associated with a which rule. we can notice that a conflicting segment  $s_5$ , which points out a conflict, is related to a rule set consisting of three conflicting rules  $r_3$ ,  $r_4$  and  $r_5$  (highlighted with a horizontal red rectangle), and a rule  $r_3$  is involved in three segments  $s_5$ ,  $s_6$  and  $s_7$  (highlighted with a vertical red rectangle). Our representation provides a better understanding of policy anomalies to system administrators with an overall view of related segments and rules.

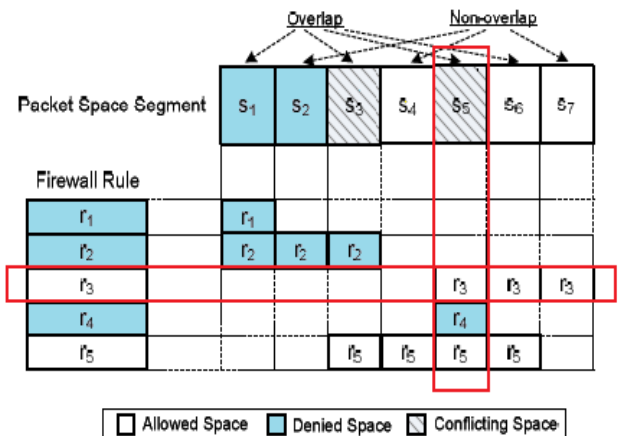


Figure 3: Representation of Policy Anomaly

## 4. Mathematical Model

Input:  $S = \{FW1, FW2, r, M(r_i), R(r_i)\}$

Where,

$S$  is a set,

$FW1$  indicates Firewall 1,

$FW2$  indicates Firewall 2,

$r$  indicates rules,

$M(r_i)$  indicates set of packets match with rule  $r_i$

$R(r_i)$  indicates set of packets match with set  $r_i$  but not match with  $r_j$  above  $r_i$  where  $j < i$ .

#### Process:

#### Privacy inter-firewall redundancy Removal

Convert each firewall to an equivalent sequence of non-overlapping rules.

$$M(nr) = R(nr)$$

Where, nr = non overlapping rules.

Firstly, this paper explains the privacy-preserving protocol for comparing a number and a range. To make sure whether a number from FW2 is in between the range [a', b'] from FW1, where, a' and b' are number from FW1, use a method similar to the prefix membership verification.

**Prefix conversion**

In this convert [a', b'] to a minimum number of prefixes, denoted as S([a', b']), whose union corresponds to [a', b']. For instance, S([11, 15])={1011, 11\*\*}.

**Prefix family construction**

This generates entire the prefixes which contains a with a itself. This set of prefixes is known as the prefix family of number a, denoted as F(a).

Let bit length of a is equal to k. The prefix family F(a) consists of k +1 prefixes where the i<sup>th</sup> prefix is obtained by replacing the last i-1 bits of a by \*. For instance, as the binary representation of 12 is 1100, then F(12)={1100, 110\*, 11\*\*, 1\*\*\*, \*\*\*\*}. It is not difficult to prove that a ∈ [a', b'] if and only if F(a) ∩ S([a', b']) ≠ ∅.

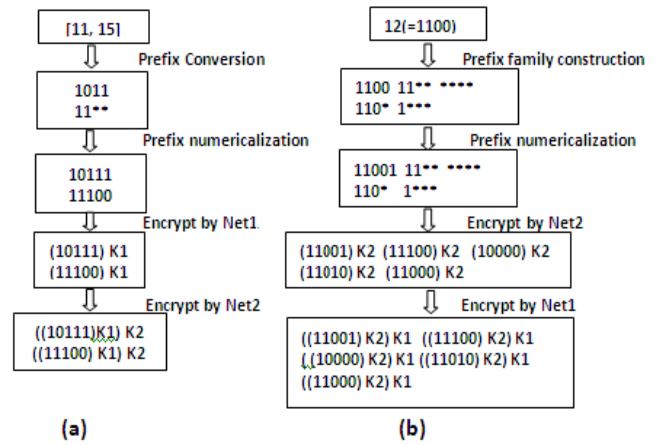
**Prefix numericalization**

This converts the prefixes obtained in the previous steps to existing numbers such that one can encrypt them in the next step. For this the prefix numericalization technique is used [19].

Given a prefix b<sub>1</sub>b<sub>2</sub> . . . b<sub>k</sub>\*. . . \* of w bits, first insert 1 after b<sub>k</sub>. The bit 1 represent a divider (saperator) between b<sub>1</sub>b<sub>2</sub> . . . b<sub>k</sub> and \* . . . \*. Then replace every \* by 0. For instance, 11\*\* is converted to 11100. If the prefix does not contain \*s, place 1 at the end of the prefix. For example, 1100 is converted to 11001.

**Comparison**

This checks whether a ∈ [a', b'] by checking whether F(a) ∩ S([a', b']) ≠ ∅, which boils down to checking whether two numbers are equal. To do this checking use commutative encryption in a privacy-preserving method. Specified a number x and two encryption keys K1 and K2, a commutative encryption is a function that satisfies the condition ((x)K1 )K2 = ((x)K2 )K1 , i.e., encryption with key K1 first and then K2 is equivalent to encryption with key K2 first and then K1.



**Figure 4:** Prefix membership verification.

**5. Experimental Setup**

**Input:**

Set of Packets: Dataset consist of packets. Following table I shows the experimental setup of set of packets which are used for practical analysis.

- SIP-** Source Internet Protocol
- DIP-** Destination Internet Protocol
- SP-** Source Port
- DP-** Destination Port

**Table 1:** Packet Pattern

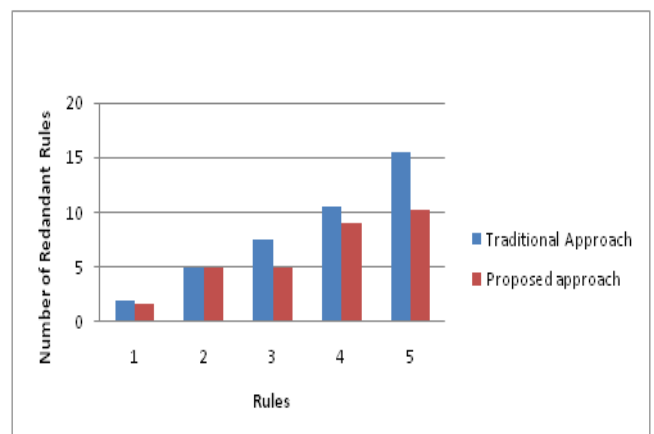
Sr. No.	Packet				
	SIP	DIP	SP	DP	Protocol
1	1.1.139.239	1.1.236.8	22	32	TCP
2	1.1.139.143	1.1.236.8	6077	3923	UDP

Set of Rules: Following table shows the example of set of rules which are used for practical analysis

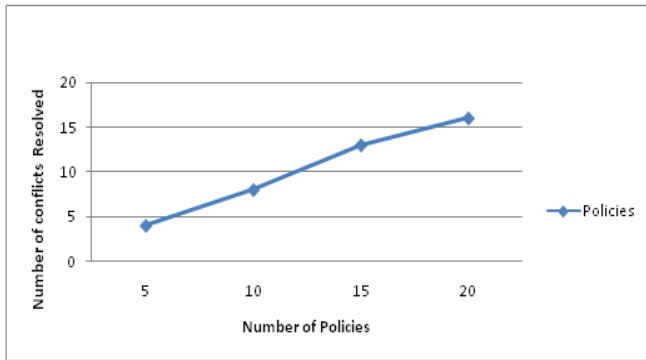
**Table 2:** Rules

Sr. No	Packet					Action
	SIP	DIP	SP	DP	Protocol	
R1	1.1.139.*	1.1.236.*	22	*	TCP	Accept
R2	1.1.139.143	1.1.*.*	*	*	UDP	Discard

**6. Result**



**Figure 5:** Number of redundant rule



**Figure 6:** Average Conflict Detection and Resolution

## 7. Conclusion

The paper represents two important mechanisms like two Cooperative firewall Optimization and Firewall security. As a result it is analyzed that the security and optimization issue are resolved effectively in the paper. The Cooperative Firewall is used to remove redundancies between two adjacent firewall policies with protecting privacy of policies. And also firewall anomaly detection techniques make secure each firewall.

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