Secure Cloud Storage with Anonymous Authentication using Decentralized Access Control

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Abstract: Security and privacy protection in clouds are challenging. Efficient search in clouds is also an important concern. A new decentralized access control scheme that supports anonymous authentication for secure data storage is proposed. In this scheme without knowing the user’s identity before storing the data the cloud verifies the authenticity. This scheme also helps in creation, modification and reading data stored in cloud and also prevents replay attacks.

Keywords: Access control, authentication, attribute-based encryption, cloud storage.

1. Introduction

Cloud computing is receiving a lot of challenges in both academic and industrial fields. Here users outsource their storage and computation to servers called clouds by the use of Internet. This frees users from the hardness of managing resources on-site. Certain services of cloud like applications (e.g., Google Apps, Microsoft online), infrastructures (e.g., Amazon’s EC2, Eucalyptus) and platforms that help developers to write applications.

Most of the data stored in cloud is sensitive like medical records and social networks. Thus, security and privacy are two major issues in cloud computing. In one hand user must authenticate itself before any transaction and on other hand cloud does not interfere with the data outsourced.

Efficient search on encrypted data is also an important issue in clouds. The query should not be known by the cloud but should be able to return the records that satisfy the query. These can be achieved by various searchable encryption methods [1] and [2]. Here the keywords sent to the cloud are encrypted and returns the results without knowing the actual keyword for the search. The problem here is that the records should have the keywords associated with them to enable search. Exact records are returned only when searched with correct keywords.

Access control is important in because only authorized users can access data. Another important fact the information must be from a reliable source. Since large data stored in cloud is highly sensitive, care must be taken to ensure access control of this sensitive information. Access control is also gaining much attention in online social networking sites where users use their personal information like pictures, videos, audios and documents.

Anonymity of the user is also necessary, if the user doesn’t need to disclose his/her identity. But the user need to prove that the information stored is from a genuine user. Apart from various cryptographic protocols like ring signatures, mesh signatures[3], group signatures [4], a new protocol called Attribute-Based signature (ABS) has been applied [4].

2. Related Work

A protocol called Attribute-Based Encryption (ABE) [7] is used that contain a set of attributes in addition to its unique ID. Two classes of ABE are key-policy ABE and Ciphertext-policy ABE. In key-policy the sender has an access policy to encrypt data. In Ciphertext-policy receiver has the access policy in the form of a tree in which attributes are leaves and have access structures with AND, OR and other gates. All the approaches takes a centralized format and allow only one KDC with a single point failure.

3. Background

In this section, a cloud storage model and the assumptions that have been made in this paper. Table 1 consists of certain notations used in this paper.

3.1 Assumptions

The following assumptions are made in this paper:
1. The cloud is honest-but-curious, cloud administrators can view the contents but cannot update it. This is a correct assumptions [5] and [6].
2. Users can read or write or both accesses to a file stored in the cloud.
3. All the communications between the users in cloud are secured by secure shell protocol (SSH).

3.2 Format of Access Policies

The access policies can be 1) Boolean functions of attributes 2) linear secret sharing scheme (LSSS) matrix, 3) monotone span programs. An example of a Boolean function is \((b_1 \land b_2 \land b_3) \lor (b_4 \land b_5)) \land (b_6 \lor b_7))\), where \(a_1, a_2, \ldots\) are attributes.

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The cloud storage model is decentralized such that there can be several KDCs for key management. Access control and authentication schemes are both collusion resistant. The proposed protocol also supports multiple read and writes on the data stored in the cloud. The costs of decentralized approaches are not individually authorized. Revoked users cannot access data or authenticate themselves, even though they are honest. A trustee can be someone like the federal government who manages social insurance numbers etc. On presenting her id (like health/social insurance number), the trustee gives her a token. There are multiple KDCs (here referred to as Aj), which can be scattered. For example, these can be servers presenting her id (like health/social insurance number), the trustee gives her a token. There are multiple KDCs (here referred to as Aj), which can be scattered. For example, these can be servers. The creator decides on a claim policy Y, to prove access to the data stored in the cloud. The costs of decentralized approaches are not individually authorized. Revoked users cannot access data or authenticate themselves, even though they are honest. A trustee can be someone like the federal government who manages social insurance numbers etc. On presenting her id (like health/social insurance number), the trustee gives her a token. There are multiple KDCs (here referred to as Aj), which are keys for signing. The secret key of KDC Aj is

\[ SK[j] = \{ \sigma_i, \omega_i, i \in H_j \} \] (1)

The public key of KDC Aj is

\[ PK[j] = \{ m(r, r)^{\sigma_i}, r^{\omega_i}, i \in H_j \} \] (2)

2) **Key generation and distribution by KDCs:** User \( U_u \) receives a set of attributes \( I[j, u] \) from KDC Aj, and corresponding secret key \( sk_{i,u} \) for each \( i \in I[j, u] \).

\[ sk_{i,u} = r^\sigma_i HF(u)^{\omega_i} \] (3)

where \( \sigma_i, \omega_i \in SK[j] \). Note that all keys are delivered to the user securely using the user’s public key, such that only that user can decrypt it using its secret key.

**Encryption by sender:** Sender decides about the access tree. Sender encrypts message M as follows:

1. Choose a random seed \( s \in Z_p \) and a random vector \( v \in Z_p^h \), with \( s \) as its first entry; \( h \) is the number of leaves in the access tree.
2. Calculate \( \beta_x = R_{x} \cdot v \), where \( R_{x} \) is a row of R.
3. Choose a random vector \( \delta \in Z_p^h \) with 0 as the first entry.
4. Calculate \( \delta_x = R_{x} \cdot \delta \).
5. For each row \( R_{x} \) of R, choose a random \( \rho \in Z_p \).
6. The following parameters are calculated:

\[ C_0 = M \cdot m(r, r)^s \]
\[ C_{1,x} = m(r, r)^{6x} \cdot m(r, r)^{\sigma_{i}(x)x} \]
\[ C_{2,x} = r^{\rho x} \]

The access policy decides who can access the data stored in the cloud. The costs of decentralized approaches are not individually authorized. Revoked users cannot access data or authenticate themselves, even though they are honest. A trustee can be someone like the federal government who manages social insurance numbers etc. On presenting her id (like health/social insurance number), the trustee gives her a token. There are multiple KDCs (here referred to as Aj), which are keys for signing. The secret key of KDC Aj is

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\[ C_{2,x} = r^{\rho x} \cdot \forall x \] (4)
\[ C_{3x} = r^{\pi(x)}x r^{\delta x} \forall x \]

where \( \pi(x) \) is mapping from \( R_x \) to the attribute \( i \) that is located at the corresponding leaf of the access tree.

7. The ciphertext \( C \) is sent by the sender:

\[ C = R, \pi, C_0, \{ C_{1x}, C_{2x}, C_{3x} \} \forall x \tag{5} \]

**Decryption by receiver:** Receiver \( U_u \) takes as input ciphertext \( CT \), secret keys \( \{ s_k_i, u \} \), group \( G_1 \), and outputs message \( M \). It obtains the access matrix \( R \) and mapping \( \pi \) from \( CT \). It then executes the following steps:

1) Calculates the set of attributes \( \{ \pi(x): x \in X \} \cap \mathcal{I}_u \) that are common to itself and the access matrix. \( X \) is the set of rows of \( R \).

2) For each of these attributes, it checks if there is a subset \( X' \) of rows of \( R \), such that the vector \( (1,0,0) \) is their linear combination. If not, decryption is impossible.

3) Decryption proceeds as follows:

   a. For each \( x \) \( \text{dec}(x) = \frac{C_{1x} e(H(u), C_{3x})}{e(x_0(x) o C_{2x})} \)

   b. \( U_u \) computes \( M = C_0 / \prod_{x \in \mathcal{I}_u} \text{dec}(x) \).

**References**


**Author Profile**

**Jessy Cherian** received the B.E. degree in Computer Science and Engineering from Mahabarathi Engineering College, Chinnasalem at 2012 and pursuing M.Tech in Computer Science and Engineering in Mount Zion College of Engineering, Kadammanitta.