

Polling, Interrupts & μ COS-II: A Comparative Timing Response Simulation Model for Wireless Processor-to-Processor Communication

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Abstract: *The paper aims to analyze the timing performance characteristics of systems employing polling, interrupts and a Real Time Operating System (RTOS). The RTOS used, specifically, is μ COS-II. The analysis incorporates a model of an Industrial system made up of microcontrollers which communicate using Zigbee protocol and data is sent and received over serial ports of the microcontrollers. One microcontroller functions as the sensor node and the other is the master which receives values from sensor node and sends out appropriate control commands in response to received sensor data. The performance of the master node is taken into account since all monitoring and control is being dedicated to this master node. A physical condition of the industrial environment is sent, and in the model we shall send real time temperature values from sensor node to master node such that when set points are exceeded then master controller sends a command to switch ON cooling system to sensor node. The time taken from when sensor node sends an out of range temperature value up to the time it will receive control command from master node is used to determine the time response of the system under the three scenarios of polling, interrupts and an RTOS. Thus the analysis aims to see how polling the serial port, serial interrupts and an RTOS affect the response, predictability and deterministic behavior of the system.*

Keywords: Timing response, Interrupts, Polling, RTOS, Predictability, Zigbee, Priorities, Task Synchronization

1. Introduction

The advances in technology have seen huge strides in how the industrial setup looks like. Traditionally we have seen the use of cables for connecting sensor and actuating devices to controllers. However by incorporating a Zigbee wireless communication model whereby the Xbee communication hardware is connected serially, system response is of prime importance. The time taken for the data received on the serial port to be processed and appropriate response taken must be taken into account. With this in mind a wireless sensor node has been designed using PIC16F877A which reads temperature values and also contains the interface to the actuators. Another node has been made which is the master node using the LPC2148, which is an ARM based microcontroller. The master node should respond to received temperature data and give out the appropriate response. An analysis is also done on how the response characteristic looks like as the number of tasks to be executed by the master node increases, that is, how this node will handle the serial available data in an environment comprising more than one task.

2. Existing Methods

From the searches done it seems not so much analysis has been done on the comparison of these three scenarios. Much attention has been put on the comparison between interrupts and polling with a main focus being towards their differences and the advantages/ disadvantages of each and not the in depth timing analysis of the two. Thus we shall give a brief overview of these forms of analysis as they are not directly linked to the study being done in this piece of work. In [2] a brief overview of differences between

interrupts and polling and advantages of each is given. Reference is made to the MSP430G2xxx microcontroller and makes use of port, timer and serial interrupts. The script also specifies how interrupts can be useful in low power applications for waking processor from idle mode. The work in [1] brings out a point that traditional models of interrupt management have an incapacity to incorporate reliability and temporal determinism which is demanded on real time systems. The author proposes a model to integrate interrupts and task handling. An analysis based on scheduling is done to evaluate and distinguish the circumstances under which the integrated model improves the traditional model.

The system in [6] addresses issues of overheads imposed by interrupts on tasks. It illustrates how interrupts, with their small execution times, reduce the utilization of the CPU if they are handled at task level. The execution time of interrupts is considered significantly less than the granularity of the online dispatcher. The model assumed is based on fixed release times and fixed deadlines. A task which has its release time fulfilled can only be delayed by higher priority task of interrupts. For analysis, interrupts have been considered as high priority tasks and included in response time analysis. The discussion which follows will now take an in depth analysis of the proposed model and how simulations are carried out to determine time response as specified in the abstract.

3. System Model

As stated earlier on the master node consists the LPC2148 while the sensor node is controlled by the PIC16F877A. The articles in [3] and [4] are earlier works by this author which illustrate the concepts of processor to processor

communication. The author has borrowed these concepts and used them in this article to establish the communication between the PIC microcontroller and the LPC2148 microcontroller. The diagrammatic setup of the two systems is shown in Fig 1 and Fig 2.

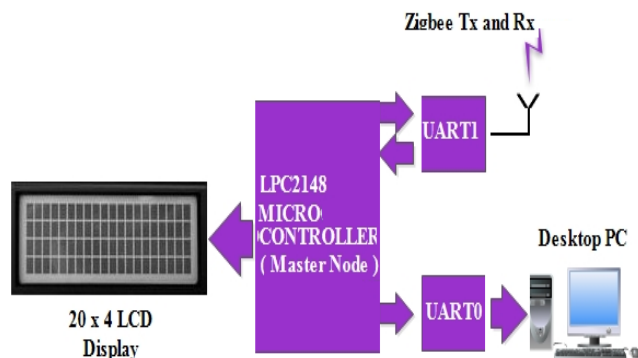


Figure 1: Master node comprising LPC2148

A. Hardware Description for LPC2148 station.

The hardware description will also illustrate how each of the hardware components contributes to system functionality. The LPC2148 is the central point to the functionality of the Master node.

1) **LPC2148:** This is a microcontroller with an ARM7TDMI based processor. The LPC series is manufactured by NXP Phillips Semiconductors.

2) **Xbee Module:** These are the hardware components used in data transmission and they use the Zigbee communication architecture. Zigbee is an IEEE802.15.4 communication interface.

3) **LCD:** The 20x4 LCD is used for display purposes and it is connected close to the master node microcontroller. It shows the updated temperature values all the time. In the timing simulations, temperature display on the LCD has been taken as one of the tasks which can utilize CPU time. In some cases the task is omitted so as to present different number of available tasks to the LPC2148.

4) **UARTS:** These are Universal Asynchronous Receiver/Transmitter modules and are used for serial data transmission. UART0 is connected to the PC for the purposes of remote temperature display. However, just like the LCD, the PC display task is not included in all the simulations. Varied numbers of tasks have been used so in some cases some tasks are exempted from the simulation. UART1 is used to connect the Xbee modules for the purpose of wireless data transmission between the Master Node and Sensor Node. As such, another Xbee module is connected to the serial port of PIC16F877A (Sensor Node).

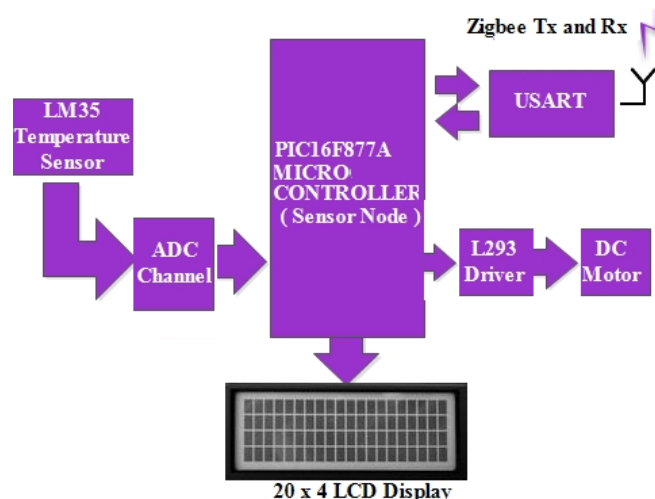


Figure 2: Sensor node comprising PIC16F877A

B. Hardware Description for PIC16F877A station

1) **PIC16F877A:** This is a medium range, 8 bit microcontroller from Microchip Company. It has been used here to coordinate the activities at the sensor node. It reads temperature values, processes them and sends the values to Master Node accordingly. It also contains the actuator interface comprising the DC motor and associated driver. Thus code to manipulate read temperature and to calculate elapsed time between sending an out of range temperature value and actuation is resident here.

2) **20x4 LCD:** At the sensor node the LCD has mainly been used for simulation purposes as a display for the number of overflows accumulated and the residue values of the Timer1 registers. These values are in decimal notation. By manipulation of these values the elapsed time can be easily calculated.

3) **LM35 Sensor:** This is a temperature sensor and its output is analogue in nature. Therefore the output has to be converted to digital format and this is achieved by connecting the sensor to an Analogue to Digital Converter channel. The LM35 gives a linear output and output voltage increases by 10mV for every 1°C rise in temperature and it can measure temperatures within the range of -55 to +150°C.

4) **USART:** The Universal Synchronous/Asynchronous Receiver and Transmitter is the serial module of the PIC16F877A and is used to interconnect the Xbee module for wireless communication with the Master Node. Some serial interrupt is used to respond to the commands send by the Master Node

5) **L293D:** The module is deigned to drive a variety of inductive loads. Examples include DC motor, relays, solenoids, bipolar stepper motors and other high current/high voltage loads in positive supply applications. It functions as a quadruple high current half H-driver. It is used in this system to drive the DC motor.

4. Method used to Determine Elapsed Time

The method used to determine time elapsed is such that whenever the PIC microcontroller sends data it will immediately start a timer. Timer 1, which is a 16-bit timer, is used and is reset to 0x0000. Timer 1 can count to 65 535 and overflow to 0 under no other advanced settings. So if the LPC2148 receives data and calculates that temperature is above set point it immediately sends a command to switch ON motor and when PIC receives this command it immediately stops timer. PIC microcontroller is configured for USART receive interrupt such that the first task inside the ISR is to stop timer. The PIC microcontroller makes calculations and number of overflows incurred and the residues of Timer Registers are used to calculate elapsed time, i.e. the number of overflows incurred and the final values of TMR1H and TMR1L registers is displayed on the 20x4 LCD and these can be used to calculate elapsed time. The delay capturing process is illustrated in Fig 3.

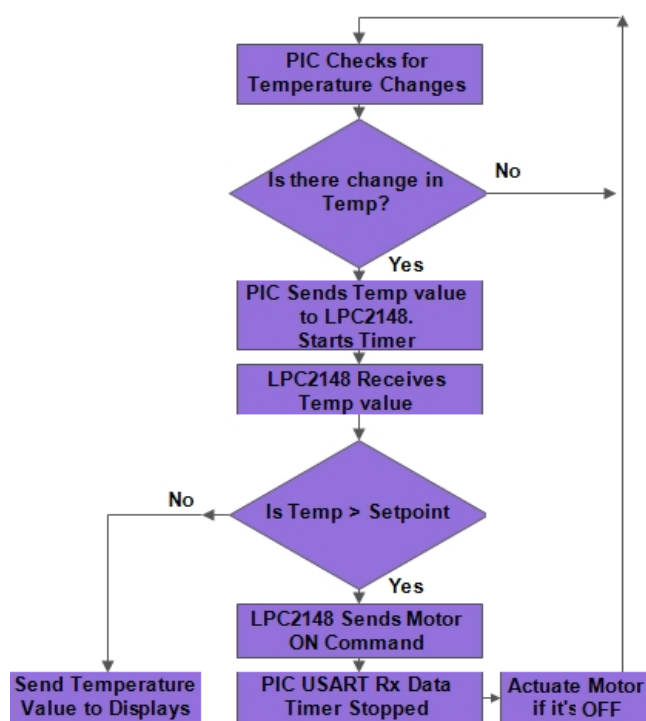


Figure 3: Diagrammatic representation of the sequence of events used in calculating elapsed time

5. System Functionality

A. Polling

Polling method utilizes a method whereby some flag is continually checked to see if it has satisfied some condition. In this setup the serial receive flag (RDR: Receive Data Ready) is continually checked to see if it is set to 1. If it is set to 1 it signifies that the Receive Buffer has some data. Thus here, tasks are executed sequentially in some specific order with the polling being one of the tasks. This order is consistent given that interrupts are disabled. We shall make an analysis on UART1 receive flag polling where different numbers of tasks are to be executed.

1) Three tasks (Polling, LCD Display, PC display via UART0): There are 3 tasks inside infinite loop and a system poll receives flag and runs the other tasks continuously. First task checks if serial receive flag has been set to signify presence of data on serial port (polling function). Second task is for displaying temperature value on LCD. Third task is to send temperature value to UART0 for display on PC. Thus serial data can be available whilst controller is still executing any one of the 3 tasks.

2) Two Tasks (Polling + PC Display via UART0): Inside an infinite loop we have a task checking/polling to see if the receive flag has been set. Another task displays temperature on PC via UART0. So system sequentially goes through these 2 tasks continually. If receive flag indicates serial data then system goes to receive the data and checks received value against set points and performs necessary action. System then returns to while (1) loop and displays temperature value received. Thus it can be seen that the serial flag can be set while system is running the temperature display routine and this will present a delay in serial response.

3) Two Tasks (Polling + LCD Display): Inside an infinite loop we have a task checking/polling to see if the receive flag has been set. Another task displays temperature on LCD. So system sequentially goes through these 2 tasks continually. If receive flag indicates serial data then system goes to receive the data and checks received value against set points and performs necessary action. System then returns to while (1) loop and displays temperature value received. Thus it can be seen that the serial flag can be set while system is running the temperature display routine and this will present a delay in serial response.

4) Single Task (Polling): Single Task inside continuous loop which polls to see if the UART1 receive flag, LSR<0>, has been set. If flag is set then a serial receive routine is initiated. The received value is then checked to see if it's not outside set limits. If it's inside limits then received data is just displayed on LCD and PC. If value is outside limit, then a command is sent via UART1 to switch ON motor driving cooling system. The value is then displayed on PC and LCD and a warning message that temperature is above limit.

B. Interrupts:

This analysis shall check the system response to UART1 interrupt. Again varied numbers of tasks are used in the implementation and simulation.

1) Two Tasks (LCD Display + PC Display via UART0): Two tasks are running inside the infinite loop and the system should run ISR task when interrupt occurs.

2) Single Task (PC Display via UART0): Apart from waiting for the interrupt the while (1) loop also runs the task which displays temperature value on PC.

3) Single Task (LCD Display): Apart from waiting for the interrupt the while (1) loop also runs the task which displays temperature value on LCD.

4) *No Task*: Continuous loop with no task running. An ISR is initiated by the presence of data on UART1. The ISR will point to the single task for receiving serial data, responding to temperature value received and display data.

C. RTOS

The RTOS, uCOS-II is a pre-emptive kernel which is portable, ROMable and scalable. It has 64 priority levels of which the highest four and lowest four are reserved for system tasks. Thus a total of 56 user tasks can be created and assigned priorities. The lower the number the higher the priority.

The analysis uses three tasks which are assigned different priorities. The three tasks are the UART receive task, LCD Display task and the PC Display task. The focus of the analysis is on the UART receive task which shall be given different priorities with the time response being noted for each case. (It should be noted that the serial task can also be responded to using interrupts in this situation but the analysis has limited the task to priority based response). The priorities shall be assigned according to the following criteria.

a) Case 1: Priority Assignments

UART Receive – Priority 5
LCD Display – Priority 6
PC display – Priority 8

b) Case 2: Priority Assignments

UART Receive – Priority 6
LCD Display – Priority 5
PC display – Priority 8

c) Case 1: Priority Assignments

UART Receive – Priority 8
LCD Display – Priority 5
PC display – Priority 6

The three cases are simulated in hardware and the timing response obtained is shown in Figure 8.

Task Synchronization and task communication

An inter task communication model can be used also in a scenario like this one since data send to LCD and PC is dependent on data received over the serial port. Task communication objects include semaphores, mutex, message queues, event flags and many more. In this case we will consider a scenario whereby the display tasks will run only if data has been received serially. Thus the UART receive task should signal the other tasks to run only if it has received some data. This concept shall be illustrated using Semaphores. In uCOS-II a semaphore is considered as an OS event and should be declared and created before being used. Consider a semaphore called *SemOne*. The Figure in 4 illustrates how we can use this semaphore.

Semaphore Operations

```
Declaring the semaphore: OS_EVENT *SemOne;
Creating the Semaphore: SemOne = OSSemCreate (0);
/* Value in brackets is the initial value of the semaphore. It can be a 1
or 0 for binary semaphores or any value from 0 onwards for counting
semaphore */
Waiting on a Semaphore: OSSemPend(SemOne, 0, &err);
/* where 1st argument is the semaphore name, 2nd argument signifies
the time out period for which task should wait for semaphore, if
Timeout = 0 it signifies infinite wait. 3rd argument is a pointer to a
variable which is used to point to error code. OS_NO_ERR means
semaphore is available*/
Posting a Semaphore: err = OSSemPost (SemOne); /*
this will increment semaphore value. Returns OS_NO_ERR if
semaphore was successfully signaled or OS_SEM_AVF if semaphore
value overflowed */
```

Figure 4: An overview of Semaphore operations

The model in Figure 5 shows the communication and synchronization of the three tasks; UART receive, LDC Display and PC Display

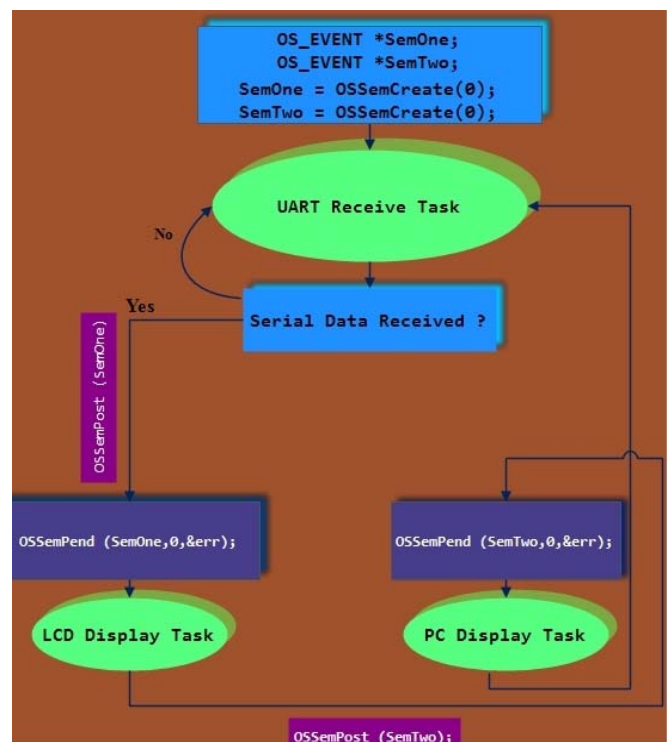


Figure 5: Task Synchronization and communication using semaphores

From Figure 5 it can be seen that unless serial data has been received no other task will execute. LCD Display task has to wait for SemOne to be available and the semaphore is only available upon being posted by serial receive task after receiving data. PC display task also waits for the semaphore, SemTwo which is posted by the LCD display task.

6. Simulations and Results

Simulations were carried out for the three scenarios of polling, interrupts and RTOS use. For each scenario varied numbers of tasks were used so as to observe system response in the presence of differing number of tasks. The diagrams

which follow show the simulation and timing analysis results. Figures 6, 7 and 8 show the timing simulation results for polling, interrupts and RTOS respectively.

6.1 Polling Response

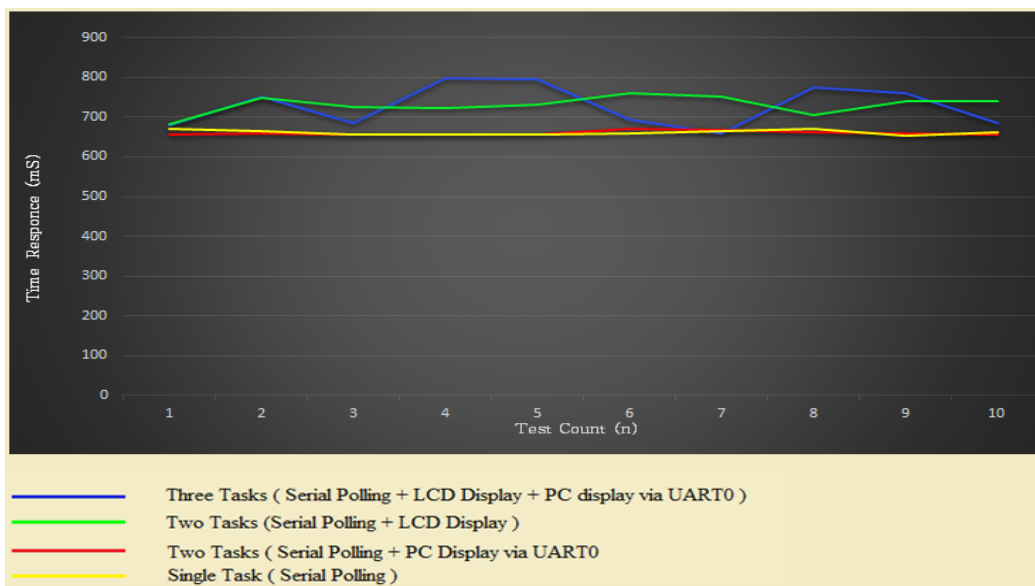


Figure 6: Timing behavior for polling response for varied number of tasks

Observation:

As the number of tasks increases the unpredictability of the system increases as shown in the diagram of Figure 6. The time taken for the system to respond to received data fluctuates over a wider range as compared to situations where the microcontroller has to poll amongst fewer tasks.

With a single task (poll task), the predictability is higher with the time response ranging between 654nS and 670nS. With three tasks, the response varies from 660nS to 798nS which is quite a wide range.

6.2 Interrupt Response

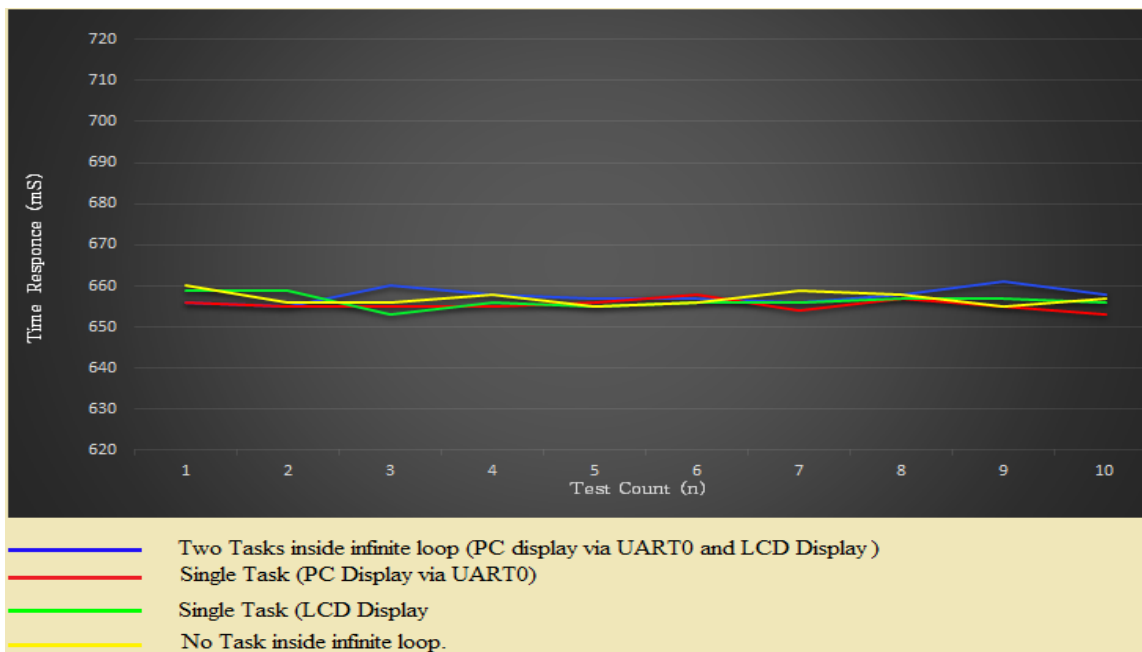


Figure 7: Timing behavior for interrupt driven response for varied number of tasks

Observation:

It can be seen that irrespective of the number of tasks the response of the system is more predictable for interrupts. The system does not take into account which task is under execution, but gives priority to the interrupt driven task and CPU leaves whatever it is doing so that it can service the interrupt signal.

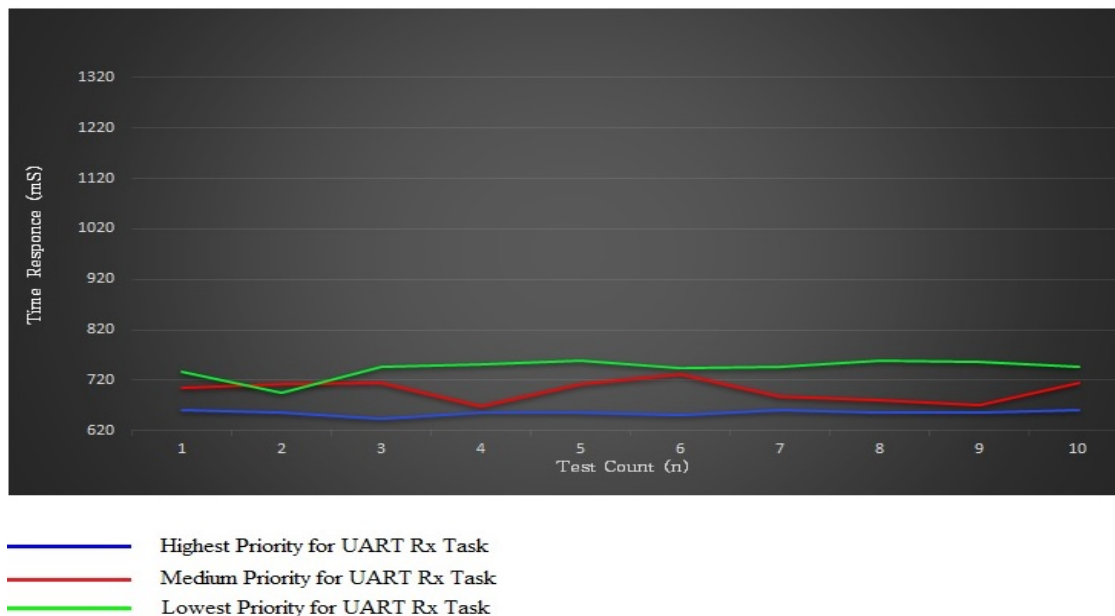
6.3 RTOS Response

Figure 8: Timing behavior for RTOS driven response for varied task priorities

Observation

In a scenario where the task being given highest priority is an external task which can also be signaled using interrupts, it can be seen that the response of the system is almost identical to that of interrupt driven response. Due to its high priority the serial receive task is responded to in a timely manner. However it has been seen that as the priority rendered to the serial task is gradually decreased the time taken to respond to serial data by the system increases and the unpredictability of the system increases as illustrated by Figure 8.

7. Overall Observation

It has been seen that for polling response time response widely varies as the number of tasks increases. Serial data becomes available while processor is executing any one of the tasks and serial data can be attended to only when the time to poll the receive flag comes. However with a reduction in the number of tasks, predictability increases and time taken to respond to serial data is reduced.

With interrupts it has been noted that irrespective of the number of tasks, serial data is responded to in the minimum possible time. The CPU is interrupted from its current task of execution and services the ISR. Therefore predictability is somehow maintained inconsiderate of task count.

The use of a Real Time OS as a measure to quickly respond to serial data has shown that the time responses of the RTOS and that of interrupts are almost identical. Thus the question which comes to mind is: What are the benefits of the RTOS? An RTOS enables the assignment of priorities to tasks irrespective of the source of the task. Not all tasks can be driven by interrupts. In such cases a task which is internal to the microcontroller can be assigned a high priority if need be. The simulation on serial response based on priority has clearly shown that the higher the priority the faster the task is responded to. Likewise any other task is given an execution slot depending on its priority.

8. Conclusions and Future Scope

This study has presented an analysis in the use of a Real Time Operating System. Basing on this analysis the author is undertaking an academic project which incorporates an RTOS ported into a master node with the purpose of monitoring and controlling an industrial setup. These timing simulations were meant to provide a concrete basis as to whether it is beneficial to use an operating system in the design. The analysis has been successfully carried out for the response to serial data. Priority assignments can serve for fast response purposes for time critical industrial processes since some processes can lead to disaster if not responded to in time. It can be noted that the use of the RTOS presents a whole lot of angles from which the analysis can be made. Simulations can be carried out to ascertain the time taken in

executing each task and make calculations to see how much a task can enter the wait state for pre-emption. This can also be used as a benchmark to make an analysis into the propagation delays incurred in the transmission of data over the wireless network. This can thus be proposed as an initiative for future work/scope into the analysis of timing simulations in microcontroller programming and communication. The study can also be extended into a situation whereby various sensor nodes are controlled by the same master node and timing analysis done to ascertain the response of the system since an industrial setup might comprise sensor nodes stationed at different locations. This will on the overall result in more time critical system performance.

experience in Computer Science and Engineering department. Her research interests include Data Mining, Computer Networks, Web Mining, Data Warehousing, Wireless Sensor networks and Microcontroller Design.

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