

Figure 3: Flow Chat of MHGA

Algorithm A:

- Step 1. Begin
- Step 2. Initialize a first part of population using Initialization heuristics algorithm
- Step 3. Initialize remaining part of population by using randomly
- Step 4. Apply Remove optimal algorithm to all tours in the initial population
- Step 5. Apply Local optimal algorithm to all tours in the initial population
- Step 6. select two parents randomly
- Step 7. Apply Crossover between parents and execute an offspring
- Step 8. Apply Remove optimal algorithm to offspring
- Step 9. Apply Local optimal algorithm to offspring
- Step 10. if offspring < any one of the parents then replace the lowest offspring of the parent
- Step 11. Shuffle any one of randomly selected from population
- Step 12. Repeat Step 3 and Step 4
- Step 13. Until number of iterations to get the optimal solution.
- Step 14. End

For example, we consider a sequence of the cities ---A ----B-C-----D-----E----F-----G-----H-----K., having distance for every city. In our scenario, Remove optimal , D is the city to be removed to perform the operations , E is the previous to the D and the next city is F, if the increases in the tour length after removing C to E to the next position , we decrease the tour length by moving into the remove city.

We can find out the decrease in the tour length: Decrease = $\text{Dist}(C,D) + \text{Dist}(D,E) - \text{Dist}(C,E)$

- I. Remove Optimal: From above Algorithm A, we defined Remove optimal to increase the tour cost of city which is badly long distance.

- a. Create a list containing the nearest N cities to a selected city.
- b. Remove optimal, removes the selected city from the tour and form a N-1 cities.
- c. Selected city which is nearest , is to reinserted in the tour and cost of the new tour length is calculated.
- d. Continue the sequence which is produces the least cost is selected
- e. Repeat 1-4 steps for each city in the tour.

- II. Local optimal: From above Algorithm A, we select n consecutive cities ($m_1, m_2, m_3, m_4, m_5, \dots, m_{n-1}$) from the tour and it arranges cities with minimum distance between the cities by searching all possible arrangements.

B. Crossover: To construct an offspring which is inheritance and adaptive the all properties from the parents structures and it acts an edge map. After getting the properties it stores information about all the connections to lead the distance between any two cities and each city have minimum two edges, maximum four edge associations from each parent

Algorithm B:

- Step 1. Begin
- Step 2. Initial city from one of the two parents from the selected city
- Step 3. Remove optimal , all occurrence of the city which is selected two parents from the edge of map
- Step 4. Selected city has entries in its edgelist go to step 5 otherwise go to step 6
- Step 5. Find city in the edgelist of the selected city and get shortest edge from the tour, and broken randomly. go to step3
- Step 6. If the parents chosen to all visited cities then stop. Otherwise randomly select an unvisited city and go to step 3
- Step 7. Selected the city with least entries in its edge list as the next selected city, choose the nearest city.
- Step 8. End.

4. Simulation Results

In genetic algorithm,. we implemented in C++ programming language for maximizing the fitness function by using $f(x) = x^2$, and the value ranges of x from 0 to 30. We implemented initial five populations of binary string and calculate x and fitness value $f(x) = x^2$. Use the tournament selection method to generate every new five populations and apply the crossover operator for generating new populations. Apply mutation operator for population to get the best fitness value.

Enter the number of Population in each iteration is: 5
 Enter the number of iteration is: 5
 Iteration 1 is:

| S.No | Population | Worst Fitness | Best Fitness |
|------|------------|---------------|--------------|
| 0 | 40269 | 0.40269 | 10.031417 |
| 1 | 1182511 | 1.82511 | 11.257072 |
| 2 | 1103802 | 1.03802 | 10.958928 |
| 3 | 1025375 | 0.25375 | 10.252110 |
| 4 | 1038920 | 0.3892 | 9.868195 |

Sum: 52.367722

Average: 10.473544
Maximum: 11.257072

Iteration 2 is:

| S.No | Population | Worst Fitness | Best Fitness |
|------|------------|---------------|--------------|
| 5 | 1153524 | 1.53524 | 8.644072 |
| 6 | 31433 | 0.31433 | 9.864642 |
| 7 | 31433 | 1.3763 | 9.045399 |
| 8 | 8313 | 0.08313 | 10.042119 |
| 9 | 1074001 | 0.74001 | 9.298878 |

Sum: 46.895107

Average: 9.379022

Maximum: 10.042119

Iteration 3 is:

| S.No | Population | Worst Fitness | Best Fitness |
|------|------------|---------------|--------------|
| 10 | 1186753 | 1.86753 | 11.619774 |
| 11 | 80292 | 0.80292 | 10.063322 |
| 12 | 158525 | 1.58525 | 9.255855 |
| 13 | 91516 | 0.91516 | 9.592331 |
| 14 | 103803 | 1.03803 | 10.959062 |

Sum: 51.490341

Average: 10.298068

Maximum: 11.619774

Iteration 4 is:

| S.No | Population | Worst Fitness | Best Fitness |
|------|------------|---------------|--------------|
| 15 | 1173828 | 1.73828 | 8.396161 |
| 16 | 64429 | 0.64429 | 10.632739 |
| 17 | 1099734 | 0.99734 | 9.900976 |
| 18 | 130327 | 1.30327 | 9.893287 |
| 19 | 1088392 | 0.88392 | 10.438623 |

Sum: 49.261787

Average: 9.852358

Maximum: 10.632739

Iteration 5 is:

| S.No | Population | Worst Fitness | Best Fitness |
|------|------------|---------------|--------------|
| 20 | 1149948 | 1.49948 | 10.060295 |
| 21 | 10010 | 0.1001 | 9.999846 |
| 22 | 130283 | 1.30283 | 9.911272 |
| 23 | 88675 | 0.88675 | 10.36997 |
| 24 | 17727 | 0.17727 | 9.883524 |

Sum: 50.224907

Average: 10.044981

Maximum: 10.36997

After the 5 Iterations, the Maximum Value is: 11.619774

Compare between Iteration 1 and Iteration 2

| Best Fitness (0-4) | Best Fitness (5-9) |
|--------------------|--------------------|
| 10.031417 | 11.619774 |
| 11.257072 | 10.063322 |
| 10.958928 | 9.255855 |
| 10.25211 | 9.592331 |
| 9.868195 | 10.959062 |

Compare between Iteration 3 and Iteration 4

| Best Fitness (10-14) | Best Fitness (15-19) |
|----------------------|----------------------|
| 10.060295 | 8.644072 |
| 9.999846 | 9.864642 |
| 9.911272 | 9.045399 |
| 10.36997 | 10.042119 |
| 9.883524 | 9.298878 |

Compare between Iteration 4 and Iteration 5

| Best Fitness (15-19) | Best Fitness (20-24) |
|----------------------|----------------------|
| 8.644072 | 8.396161 |
| 9.864642 | 10.632739 |
| 9.045399 | 9.900976 |
| 10.042119 | 9.893287 |
| 9.298878 | 10.438623 |

Compare between Iteration 5 and Iteration 1

| Best Fitness (20-24) | Best Fitness (0-4) |
|----------------------|--------------------|
| 8.396161 | 10.031417 |
| 10.632739 | 11.257072 |
| 9.900976 | 10.958928 |
| 9.893287 | 10.252110 |
| 10.438623 | 9.868195 |

Graphs Representation for Best Fitness Function

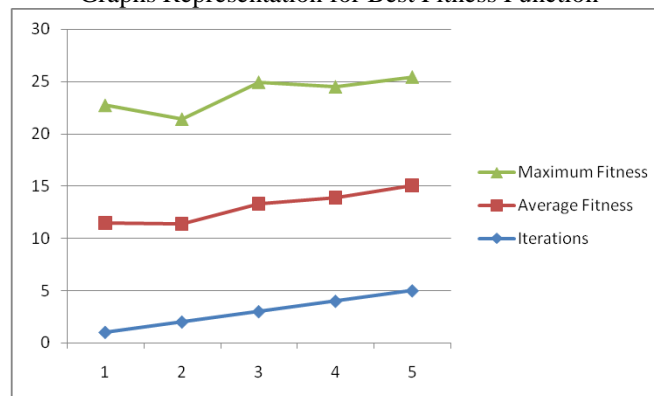


Figure 4: Best Fitness Value

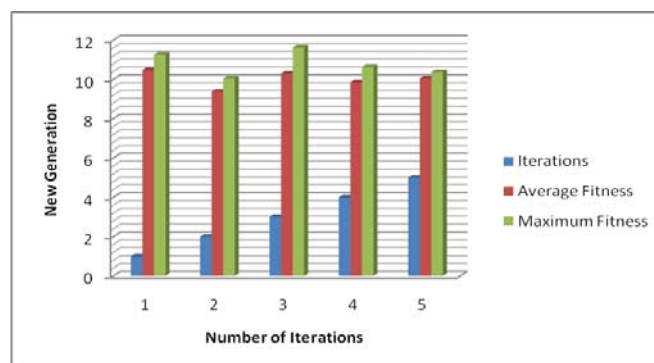


Figure 5: Compare between 5 iterations for Best Fitness Function

Compare between Number of Iterations

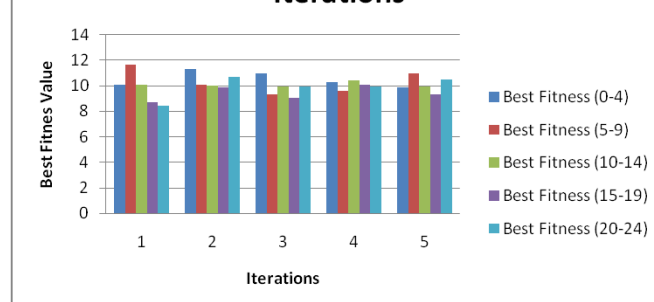


Figure 6: Comparing between Number of Iterations

5. Conclusion

In this paper, we optimize the energy consumption in datacenters in all levels. Our energy model is based on the Dynamic voltage scaling and improves the load balances by using Dynamic Load Balance Distribution Algorithm (DLBDA). We address the resource allocation and solved by adopting the features of Hybrid genetic algorithm and implemented the Modified Hybrid Genetic algorithm (MHGM) allocated the resources based on the Best Fitness values from the Populations in every iteration levels.

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