

# Survey on Interleavers in IDMA System

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**Abstract:** *In recent days, a new multiple access system called interleave division multiple access (IDMA) has provide greater attention to provide interleavers. A review of IDMA technology provided in communicating system based on different types of interleavers. This paper provides a study of IDMA technique with random interleaver, helical interleaver, tree based interleaver, prime interleaver, algebraic interleaver, block interleaver, square interleaver and power interleaver. At very low receiver complexity, IDMA has been provided a single user communication system with multiple user surroundings. Interleavers are provided by the means of user separation. Interleaver is used to improve forward error correcting performance and different signals are recognized from different user with the help of interleavers. We compare different interleavers based on memory requirement, bit error rate and computational complexity. In 4G communication systems, some basic requirements are receiver cost, ISI treatment, asynchronous control, power efficiency, user number and spectral efficiency should high.*

**Keywords:** Interleaver, bandwidth requirement, memory requirement

## 1. Introduction

There is various multiple access technique in the wireless communication system. It is one of the techniques that also used in mobile communication system. Various techniques have been studied that have to improve the bandwidth, efficiency complexity and multi-user detection (MUD) is always and there are the number of users which increases that can be accommodated within each cell. For very high mobility, the data rates are up to 100 Mbps and for low mobility, it is up to 1 Gbps are required. These requirements for the system are considered as fourth generation (4G) systems. Data rate of 3.6 to 7.2 Mbps are useful in third generation system. While in modern communication system, Code Division Multiple Access (CDMA) has greater impact in wireless communication. CDMA system is provided with high complexity and main problem in CDMA is multi user detection. It offers well known features such as dynamic channel sharing, soft capacity, reuse factor of one, low dropout rate and large coverage, ease of cellular planning, robustness to channel impairments and immunity against interference.

Interleaver division multiple access (IDMA) is a one of the technique that relies on different interleavers to separate signals from different users in a multiuser spread-spectrum communication system. In CDMA, many users are distinguished by different signature sequences, while IDMA distinguishes users by chip level interleaving methods. The main limitation with conventional CDMA systems is multiple access interference (MAI) and intersymbol interference (ISI). As per Kuldeep choudhary, P S Sharma, CDMA systems are having high degree of big problem [1]. For conventional CDMA system, the characteristic feature is the use of signature sequences for user separation. As per Shuang Wu, Xiang Chen, interleaving is generally occurred in between forward error correction (FEC) coding and spreading and is imply to the fading effect [5].

The key principle of Interleave division multiple access scheme is different interleavers for different users. If the interleavers are assigned for the user, then it obtains the

interleavers that are used in IDMA system should be efficient and least complex. In the transmitter of IDMA scheme, there is a chip level interleaver is followed by spreading process which is different from conventional CDMA scheme. Hence IDMA inherits many advantages from CDMA especially, chip by chip interleaving can be against fading and mitigation of other cell user interference. Since interleavers are used for the means of separating users in IDMA systems, it is necessary to design them properly. For random interleavers, the interleaver matrix has to be transmitted to the receiver, which can be very costly. For reducing the memory consumption and some amount of information must be exchanged between mobile stations and base stations.

The design of interleavers has significant influence on the throughput of iterative MUD receivers. But all the interleaver schemes can only be operated serially, which will have a great limit the processing throughput in hardware or software implementation. Different types of interleavers provided in IDMA communication system. The performance analysis of different interleavers such as random based interleavers, master random based interleaver, tree based interleavers and prime interleavers are compared in IDMA system in fading environment.

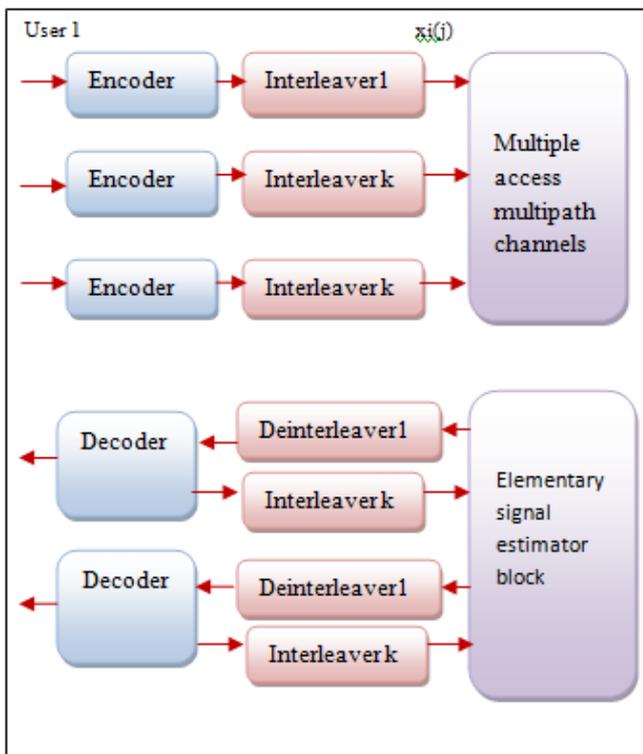
## 2. Interleave Division Multiple Access Schemes

IDMA stands for interleave division multiple access. This scheme is used for spread spectrum mobile communication systems. In IDMA scheme, number of user is differentiated by different chip-level interleavers instead of by different signatures. IDMA is one special form of CDMA in which bandwidth expansion is entirely performed by low-rate coding. Interleave division multiple access is more conveniently used system. This scheme inherits many advantages from CDMA such as cross-cell interferences, dynamic channel sharing, asynchronous transmission and robustness contrast to fading. It has very low complexity multiple user detection techniques that applicable to systems with large numbers of users in

multipath channels. Interleaving has some advantages over scrambling that are important for cell edge subscriber stations to receive broadcast services such as common signaling broadcast. Many advanced transmitting techniques used for unicast but cannot be used for broadcasting. In IDMA, data streams are separated by different interleavers rather than by different spreading codes. The data rate can be taken by superimposing interleaved data streams and many encoded streams. Whether there are other system designs, in opposite to that channel coding is an integral part of the system design. Separation of the data streams can be done at the receiver in iterative way with low complexity. These properties are more advantageous in case of multi user detection at the uplink. Hence it makes IDMA an attractive candidate for the 4G uplink.

## 2.1 IDMA Transmitter and IDMA Receiver

The block diagram of IDMA scheme is shown in "Fig. 1," for K users. The principle of iterative multi user detection (MUD) is used which is a promising technique used for multiple access problems (MAI) that are illustrated in the lower part of "Fig.1". The turbo processor which includes elementary signal estimator block (ESEB) and number of K decoders. The ESEB partially resolves MAI without considering encoding. The outputs of the ESEB are then passed to the decoder for further refinement using the encoding that is used for de interleaving block. The decoder outputs are fed back to the ESEB for improvement in estimation in the next iteration with proper user specific interleaving. This iterative procedure will be repeated for number of times and also terminated if a given certain criteria is fulfilled. After the final iteration, the decoder produces hard decisions on the information bits [4].



**Figure 1:** Transmitter and Receiver structure of IDMA system

## 3. Different Types of Interleavers

The interleaver design is play very important role in the efficiency of IDMA system. It provides decorrelation between adjacent bit sequences that are provided in the orthodox turbo coding and decoding. It also provides uncorrelate to the various users. The correlation amongst interleavers is used to measure signals that are used from other users that affecting decoding process of a specific user. If decorrelation is better, then less iteration are required for detection in IDMA multiuser detection (MUD) mechanism. The interleaver providing decorrelation is used for reducing the MAI from other user. Hence it helps in the convergence of detection process. A set of interleavers is lie under these two criteria. First, it is easy to generate from transmitter and receiver so that do not store or communicate many bits in order to steady upon an interleaver and second is that no two interleavers are in the form of set collide. It also shown that correlation between interleavers can be used to formulate a collision criterion, where zero correlation called as orthogonality which implies that no collision. In case of IDMA, the transmitter needs to transmit the interleaver matrix which consists of interleaving pattern of the users to the receiver. If greater the size of the interleaver then more bandwidth and resources are to be used. Also in IDMA, greater the size of interleaver then more the orthogonality is achieved between them. Interleaving is generally a process of rearranging the ordering of a data sequence from one to one deterministic format. Interleaving is a technique that will enhance the error correcting capability of code. Interleaving is used before the information data is encoded by the second component encoder in turbo encoder. The main role of an interleaver is constructing a long block code from small memory convolution codes. The final role of the interleaver is that break the low weight input sequences then increase the code free Hamming distance or reduce the number of code words with very small distance in the code distance spectrum. The size and structure of interleavers also play a major role in the performance of turbo coding. There are a number of interleavers, which either can be implemented.

### 3.1 Random interleaver

Random interleavers can interleave the data of different users with different pattern. These patterns of scrambling the data of users are to be generated. These scrambling of data causes burst error of the channel are randomized to the receiver end. With the help of random permutation, the user specific random interleaver is used to rearrange the elements of its input vector [8]. These incoming data can rearrange using a series of generated permuter indices. A permuter is generally a device that helps to generate the pseudo-random permutation of given memory addresses. These data is arranged in order of pseudo-random order of memory addresses. When random interleavers are work for the purpose of user separation, then it requires a lot of memory space will be at the transmitter and receiver ends for storage purpose. Hence considerable amount of bandwidth will be consumed for transmission of all these interleaver. Thus at receiver end, computational

complexity will be increases. After randomization of the burst error that has to rearrange the whole block of the data afterwards it can now be easily detected and corrected. Spreading is one of the essential characteristic of random interleavers. Random interleaver is used in IDMA system for user separation, so interleaver must satisfy certain design criteria. Interleavers of different users do not collide for user separation. The property of minimum collision amongst user specific interleavers depends on property of orthogonality. It is referred as an important factor in generating the interleavers [1].

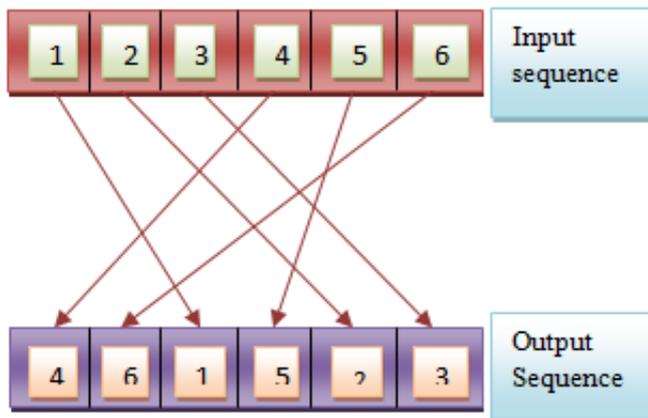


Figure 2: Random interleaver

Two interleaver are said to be orthogonal when

$$\pi(j) = \pi(i) \tag{1}$$

Whenever, correlation between two code words x and y satisfy the following relation

$$C(\pi(i), x, \pi(j), y) = \{ \pi(i) (f(x)) \pi(j) (f(b)) \} \tag{2}$$

Above condition is for checking orthogonality from “Fig.3,” between user-specific interleavers [8]. If the interleaver is not randomly generated, the system performance degrades considerably and the MUD is unable to resolve MAI problem at the receiver resulting in higher values of bit error ratio (BER) [5]. If random pattern are generated for different interleavers, then MUD resolves MAI problem. Hence it results in improving BER.

### 3.2 Helical interleaver

Helical interleaver reduces the burst error correction capability. In this interleaver, interleaver indices are read in a deterministic order for generating helical interleaver. Initially the data elements in helical interleaver are arranged row wise and column wise, then data can be read in diagonal-wise as shown in “Fig. 1,” Helical interleaved index is calculated by using the following formula given as,

$$j = i(k_x + 1) \text{mod}(k_x k_x), i = \{0, 1, 2, \dots, k_x, k_x - 1\} \tag{3}$$

Where i= original index value, j = helical index value,  $k_x =$  x dimension bits,  $k_y =$  y dimension bits [8].

The generation of helical interleaver can be explained as first to generate primary interleaver of length L, then write data sequence in matrix form having  $k_y$  rows and  $k_x$  column, where L is length of interleaver can be depicted as  $L = k_y \times k_x$ . Then, helical interleaver can be obtained with the generation of primary interleaver by reading interleaver element column wise. In final step by reading the value of interleaver indices from diagonal element of the matrix in cyclic manner, other interleaver can be easily generated.

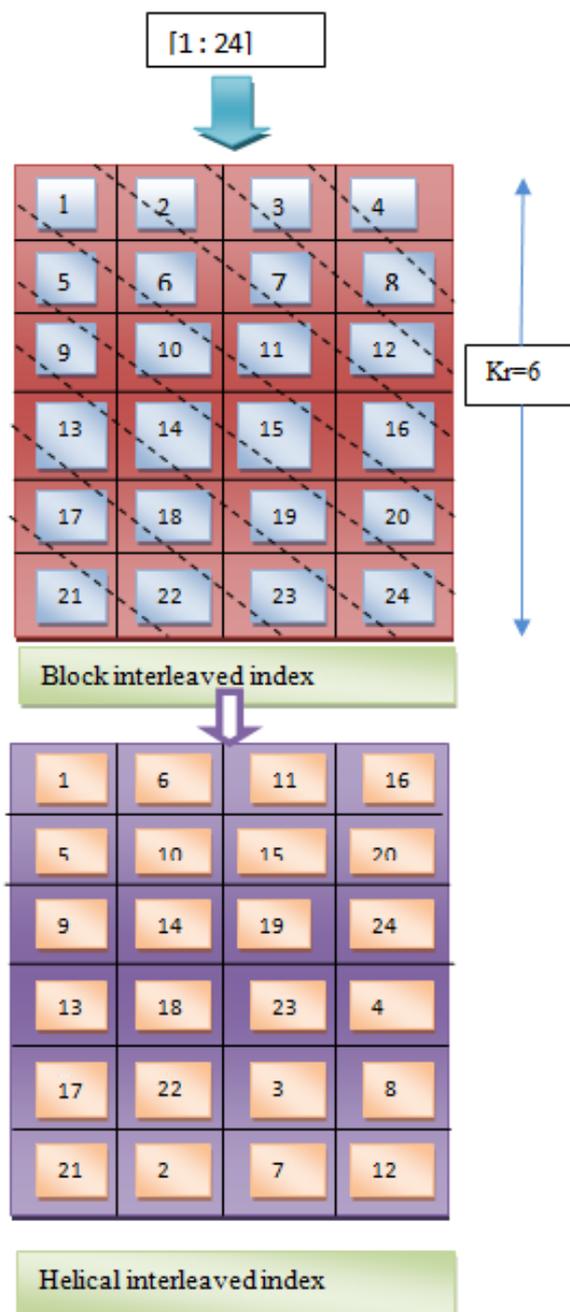


Figure 3: Helical interleaver

Mathematically,  $i$ th helical interleaver can be given as,

$$\pi_i[k] = \pi_i[l_{mod L}], 0 \leq k < k_c \tag{4}$$

Where

$$l = k_{mod k_r} \cdot k_c + \left(\frac{|k|}{|k_r|}\right) + (k_{mod k_r} \cdot (i-1)) \mod k_c \tag{5}$$

If parameter is chosen properly then interleaver element that have to used are spread out in an effectively. Different users are used by different interleavers. Constructing a layer specific shift, can be optimized as

$$\pi_i[k] = \pi(l + i, s) \mod k_c \tag{6}$$

Where  $S$  represents the interleaver shift constant.

### 3.3 Tree Based Interleaver

Tree Based Interleaver (TBI) is used as different chip-level interleaving sequences used for different users in an IDMA system. The operation involves reducing computation complexity. This method of generating interleaving sequence solves the memory cost problem and also reduces the amount of information exchange between mobile stations and base stations that required for specify the interleaver. The algorithm for TBI is based on the selection of combination of two master interleavers. The odd number of users is taken upside and even number of users is taken downside. In that way, a large number of users may be allocated with user specific interleavers having very less complexity. By using a combination of randomly selected master interleavers, interleavers are designed.

Here  $\pi_1$  and  $\pi_2$  are two master interleavers which are randomly selected. The  $\pi_2$  interleaver is reserved for initiation for lower branch. Upper branch is selected for odd user count while lower branch is selected in case of user count is even. For first user interleaver will be  $\pi_1$  while the interleaver will be  $\pi_2$  for second user. In case of third user it will be  $\pi_1(\pi_1)$  and the interleaving sequence will be  $\pi_2(\pi_1)$  for fourth sequence. The memory requirement of tree based interleaver is extremely low as compared to the random interleaver and it is slightly high in comparing with master random interleaver [5]. The IDMA system in comparison with random interleaver will imposes the problem of extra bandwidth consumption in the channel in responding to the high memory requirement at the transmitter and receiver ends. Hence the memory requirement for storing the user specific interleavers is user dependent in case of random interleavers. In IDMA scheme, it is found to be at minimum level. For tree based interleaver, the memory requirement is observed to be little bit high in comparison to that required in case of master random interleaver. While it is extremely less in compare with requirement in case of random interleaver.

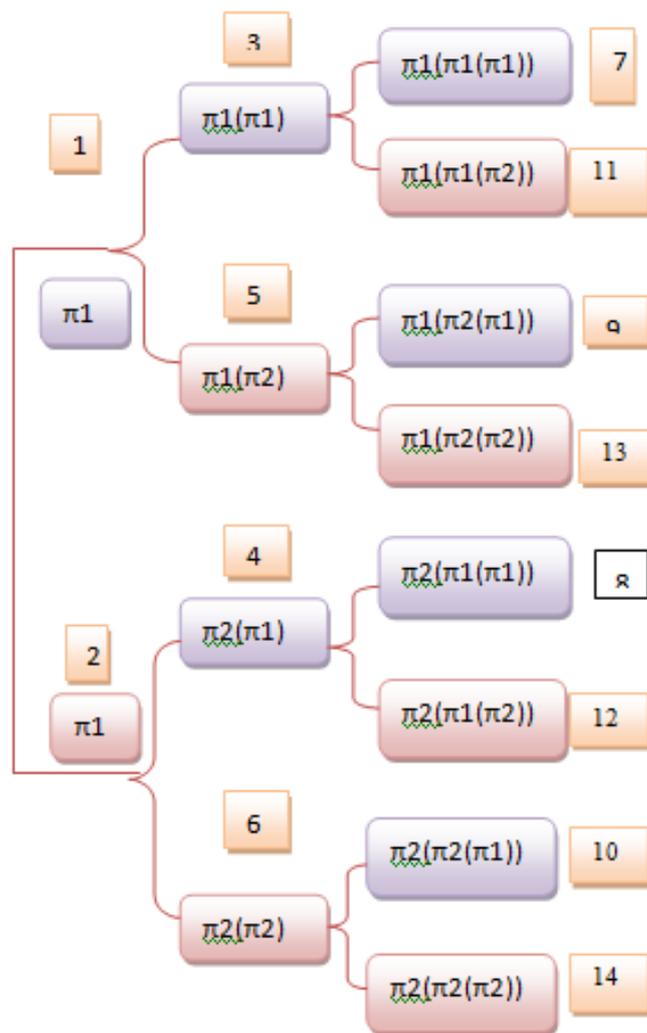


Figure 4: Tree based interleaver

The two interleaver combines in such a fashion as shown in “Fig.5” to form interleaving mask for the users. The interleaving sequences related to 14th user needs only 2 cycles,

$$\pi_{14} = \pi_2(\pi_2(\pi_2)) \tag{7}$$

$$\pi_7 = \pi_1(\pi_1(\pi_1)) \tag{8}$$

Where  $\pi_1$  and  $\pi_2$  are two master interleavers.

### 3.4 Prime interleaver

In prime interleaver, user specific interleaver is designed by using a combination of both master interleavers. The scheme is used very rarely in terms of bandwidth requirement and bit error rate. However, there is space required for development of other efficient interleavers for IDMA scheme. In IDMA scheme, different users are assigned for different interleavers that are very weakly correlated. For generation of interleavers, the computational complexity and memory requirement should be small. The Prime Interleaver is use to minimize the bandwidth and memory requirement that occur in other available interleavers with bit error rate (BER) performance are compared with the random interleaver

[5]. While in the generation of prime interleaver, the prime numbers are considered as seed of interleaver. In this scheme, user specific seeds are assigned to different users. For finding the mechanism of prime interleaver, consider a case of interleaving of  $n$  bits with seed  $p$ . First, let Gallois Field  $GF$  is  $n$ . Now, the bits are interleaved with a distance of seed over gallois field  $n$ [5]. In case, if  $\{1, 2, 3, 5, 6, 7, 8 \dots n\}$  are consecutive bits that have interleaved with seed  $p$  then location of bits after interleaving will be as follows;

- 1 => 1
- 2 =>  $(1 + p) \bmod n$
- 3 =>  $(1 + 2p) \bmod n$
- 4 =>  $(1 + 3p) \bmod n$
- ..
- $n$  =>  $(1 + (n-1)p) \bmod n$

For Example interleave 8 bits as  $\{1, 2, 3, 4, 5, 6, 7, 8\}$  and interleave these bits with seed 3 then the new location of bit will be as follows

- 1 => 1
- 2 =>  $(1 + 1*3) \bmod 8 \implies 4$
- 3 =>  $(1 + 2*3) \bmod 8 \implies 7$
- 4 =>  $(1 + 3*3) \bmod 8 \implies 2$
- 5 =>  $(1 + 4*3) \bmod 8 \implies 5$
- 6 =>  $(1 + 5*3) \bmod 8 \implies 8$
- 7 =>  $(1 + 6*3) \bmod 8 \implies 3$
- 8 =>  $(1 + 7*3) \bmod 8 \implies 6$

Here, the new order of bits will be  $\{1, 4, 7, 2, 5, 8, 3, \text{ and } 6\}$  [5]. The bandwidth requirement for the prime interleaver is smaller than other available interleavers. As very small amount of memory is required for the transmitter and receiver side, hence there is a need of seed value. While in master random interleaving scheme, the computational complexity and transmitter and receiver end is quite high. The prime interleaver results in reducing the computational complexity that occurs in master random interleaving scheme. Due to the computation of user specific interleaver, it is higher to that of tree based interleaving scheme.

### 3.4 Algebraic Interleaver

The algebraic interleaver is helps to rearranging the elements of its input vector using a permutation that can be algebraically derived. The parameter having number of elements  $N$  indicates that various numbers are available in the input vector. This will have column vector input signal. The output signal can inherit its data type from the input signal. The Type parameter indicates the algebraic method that having block users to generate the appropriate permutation table.

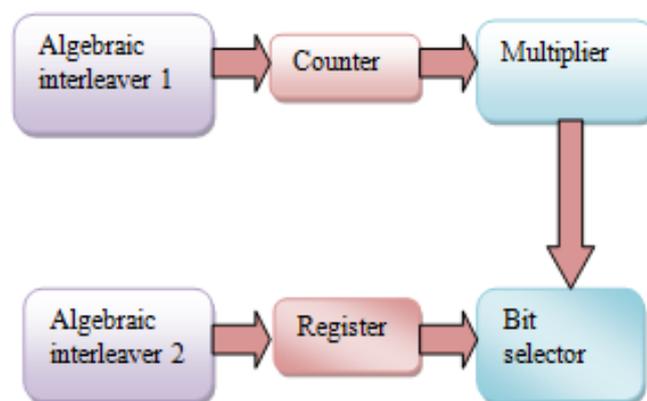


Figure 5: Algebraic interleaver

A Takeshita-Costello interleaver is uses a length- $N$  cycle vector whose  $n$ th element is

$$C(n) = \bmod(k.n(n-1)2, N) + 1, n \tag{9}$$

for integers  $n$  from 1 and  $N$ . The intermediate permutation function is obtained by using the following relationship:

$$\pi(c(n)) = c(n+1) \tag{10}$$

Where

$$n = 1: N \tag{11}$$

The permutation vector is the result of cyclic shifting of elements of the permutation vector  $\pi$ , by the cyclic shift parameter,  $h$ .

### 3.5 Power interleaver

The power interleaver  $\emptyset$  is needed to be stored in particular way. Let the power interleaver be  $\pi_1 = \emptyset$ . After completing the detection cycle for user 1, the interleaver can be taken from  $\pi_1 = \emptyset$  to  $\emptyset(\pi) = \emptyset^2$ . This will continues in recursive way. The new interleaver generation method is called as power interleavers that can take the place of random-interleavers without performance loss [3].

The drawback of this scheme is that it has higher access time for user securing  $\emptyset^n$  interleaver where  $n$  is the user number. Simulation result shows that similar results have been obtained as that achieved with Random Interleavers, but considerable amount of memory space has been saved.

### 3.5 Block interleaver

The Block Interleaver block having rearranging of the elements of its input vector without repeating of elements. If the input is having  $N$  elements. The Permutation parameter is contains column vector of length  $N$ . The column vector having indices, hence input elements are in form the length- $N$  output vector that is, Output  $(k) =$  Input (Permutation vector  $(k)$ ) for each integer  $k$  between 1 and  $N$ . The contents of permutation vector having integers between 1 and  $N$  and it have no repetitions. Both the input

and the permutation vector parameter must have column vector signals.

This block can be output sequences that vary in length. To get more information for sequences which varies in length also with variable size signals. The output signal can inherit data type from the input signal [3]. For example, If Permutation vector is [4;1;3;2] and the input vector is [40;32;59;1], then the output vector must be [1;40;59;32]. Hence all of these vectors should have same length and that the vector Permutation vector is a permutation of the vector [1:4].

#### 4. Comparison of Interleavers

The comparison between different interleavers used in IDMA technologies have been made on the basis of memory requirement, bandwidth requirement, complexity, bit error rate (BER) for random based interleaver, helical interleaver, tree based and prime interleavers.

**Table 1:** Comparative table for different interleavers

Interleavers				
Parameters	RI	HI	TBI	PI
Memory requirement	High	Low	Low	Lowest
Bandwidth requirement	1.5* 10(6)	0.01*1 0(6)	0.02* 10(6)	0.0001* 10(6)
Complexity	High	Very high	Low	Little high than TBI
Bit error rate for Eb/No=10(24 users)	10(-4)	10(-4)	0.04* 10(-4)	0.5*10(-4)
BER in coded environment for Eb/No=10(24 users)	0.6* 10(-5)	0.6* 10(-5)	0.4* 10(-6)	0.4* 10(-6)
BER in uncoded environment for Eb/No=10(24 users)	0.6* 10(-4)	0.2* 10(-4)	0.2* 10(-5)	0.2* 10(-5)
Specific user cross correlation	Low	Low	High	High

#### 5. Conclusions

The four interleavers as random, helical, tree based, and prime interleaver are compared based on the BER performance, memory requirement and complexity for their generation. Tree based and helical interleavers have quite similar BER performance compared to that of random interleaver. Also, helical interleaver has less complexity than tree based and random interleavers taken in to consideration. Thus, helical and tree based interleaver are more suitable for IDMA system. Prime Interleaver is very easy to generate and is better than other interleavers in terms of bandwidth, memory, bit error rate etc. Prime Interleaver is better than random interleaver in terms of computational complexity. With tree based interleaver, the proposed interleaver seems to be having little bit more complexity due to involvement of higher calculation for user specific interleavers. These two interleavers, tree based and helical interleaver outperforms the random interleaver in terms of bandwidth and memory

consumption, which results in saving a lot of memory space while transmission of interleaver sequences, hence it results in increasing the efficiency of IDMA system.

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