

# High Speed Advanced Encryption Standard Using Pipelining

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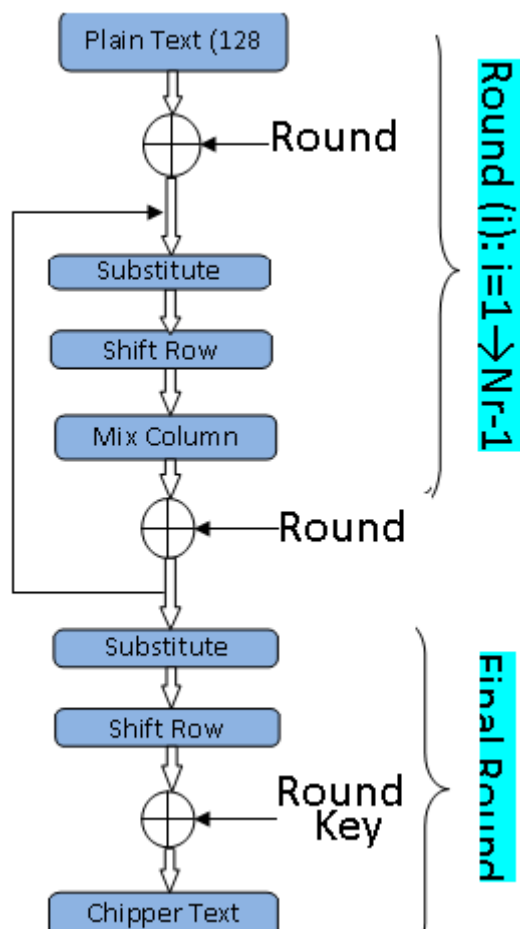
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**Abstract:** In this paper we have proposed high throughput by swapping the AES algorithm internal stages in this proposed work shift row is operated before sub bytes (substitution bytes). In this proposed operation the AES encryption operation will not effect, with this process is streamlines the processes a 4 block of data rather than 16 block. The advantage of this is we can save area. This process repeats for 10 cycles and with the help of this we can encrypt 128 bits data with higher throughput. We have evaluated this performance of higher throughput and hardware area in Xilinx's 12.2 vertex 4 XC4VFX140-11FF1517.

**Keywords:** Advanced Encryption Standard (AES), Register Transfer Logic (RTL), Very High Speed Integrated Circuit Hardware Description Language (VHDL), Galois field (GF), Byte Substitution (Sub Byte)

## 1. Introduction

Advanced Encryption Standard (AES) is used for security purpose in Military application. All Cryptographic algorithm are used for security services for various application all the encryption technique are used in government and secret military communication. The basic algorithm of AES can easily understand by the flow chart given below:



The AES used for encrypt and decrypt 128 bit plain text block. To encrypt this plain text we required 3 modes: 128

bit, 192 bit 256 bit. Each has corresponding number of round. To encrypt the data we required 128 bit matrix and each row contains four bytes of group. The 4\*4 Matrix is given below.

$$X = \begin{bmatrix} x_0 & x_4 & x_8 & x_{12} \\ x_1 & x_5 & x_9 & x_{13} \\ x_2 & x_6 & x_{10} & x_{14} \\ x_3 & x_7 & x_{11} & x_{15} \end{bmatrix}$$

Figure 1: Matrix

The AES algorithm consists of four different simple operations.

- Byte substitution (Sub Bytes)
- Shift Row
- Mix column
- Add Round Key

### 1.1 Sub Bytes:

All bytes are processed separately and it is a non linear byte substitution. Sub byte is invertible and constructed by the composition of following two transitions;

Inversion in the  $GF(2^8)$  field and modulo an irreducible polynomial is given by:

$$M(x) = x^8 + x^4 + x^3 + x + 1$$

Affine transformation defines by as following:

$$Y = AX^{-1} + b$$

Where A is 8\*8 fixed matrix and b is 8\*1 vector matrix.

Shift Row:

In this stage

- The first row is not change.
- The second row is circular shifted by 1 byte to the left.
- The third row is circular shifted by 2 byte to the left.

- The fourth row is circular shifted by 3 byte to the left.

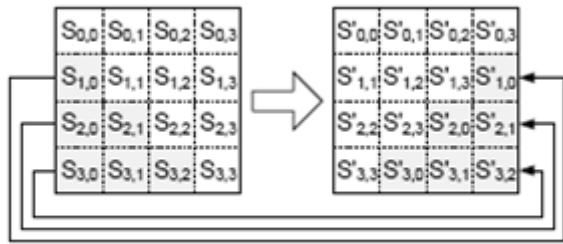


Figure 2: Shift Row

### 1.2 3 Mix column

The mix column transformation operates column by column and these columns will be consider as a four term polynomial. The column are consider as four term polynomial over  $GF(2^8)$  are multiplied  $x^4+1$  with the fixed polynomial  $a(x)$  is given by:

$$a(x) = \{03\}x^3 + \{01\}x^2 + \{01\}x + \{02\}$$

### 1.3 Add Round Key:

In this transformation the 128 bit of stage are bitwise xor with the 128 bit of the round key. The operation is viewed as a column wise operation between the 4 bytes o a state column and one word of the round key.

## 2. Methodology

The AES can be viewed as a encryption algorithm and in our proposed methodology we will do some changes in the basic AES algorithm and improve it's performance. The first change is we will swapping two block that is sub byte and shift row with this change the avalanche effect will improve and the second change we will perform is the introduction of pipeline in the basic AES algorithm with this change the throughput will increase and the area (slices) will reduce.

### 2.1 Swapping of Blocks

In the AES algorithm the blocks "sub byte" and "shift row" are swapped together but the operation of AES algorithm will not be affected when the shift row blocks is operated before the byte substitution block then the avalanche effect of the AES algorithm will be increases.

### 2.2 Avalanche Effect

It is the important characteristics for encryption algorithm. This property can be seen when changing ne bit in the plain text and then watching the change in the outcome of at least half of the bits in the cipher text. One purpose of the avalanche effect is that when changing in the one bit there is a large change then it is harder to perform an analysis of cipher text when trying to come up with an attack.

Consider a function  $F1\{i, j\}^n$  here  $\{i, j\}^n$  satisfy avalanche criteria when one input is change at least half of bit in the output bit change. Where  $i, j$  are the input and output bits. As per avalanche criteria:

Total charge in  $j^{n/2}$  avalanche variable computed over whole input size  $2^n$  in range  $0 \leq w(a_j) \leq 2^n$  from the above equation we can calculate avalanche parameter of  $i$  as :

With above formula it can be proved that probability of change the output bit when only one or  $i^{th}$  bit of the input is changed is half.

### 1.3 Pipeline

In our design we uses the pipeline architecture of the AES pipeline stages as well as it divides the AES into ten AES pipeline stages as well as it divides the AES into ten stages yielding an overall of 110 pipeline stages then resulting speed in terms of throughput rate and implementation area is evaluated and compared with existing design implementation in terms of same gate technology factor. In our design we uses register in between every block of AES. The register is behaves like memory and used to store some data for some instant of time:

### 1.4 Throughput

Throughput can be calculated by using following equation:  
**Throughput = (Total plaintext in byte encrypted)/(Total execution time)\*100**

Table 1: Throughput comparison

Existing technique	File size	%
AES	50KB	70%
DES	50KB	60%
Blowfish	50KB	38%
Caesar cipher	50KB	2%
Vigenere cipher	50KB	4%
Platfair cipher	50KB	8%

## 3. Results

In this section we will discussed about all the result come from our proposed work. In our proposed work we will give a clock cycle for all the four blocks will operate according to the clock cycle in the input. In our proposed design the number of slices will be reduce and slices utilization percentage is also reduce. The main parameters of our results are number of slices, throughput and frequency. The result come from simulation (synthesize summary) and the RTL view of proposed design is given below.

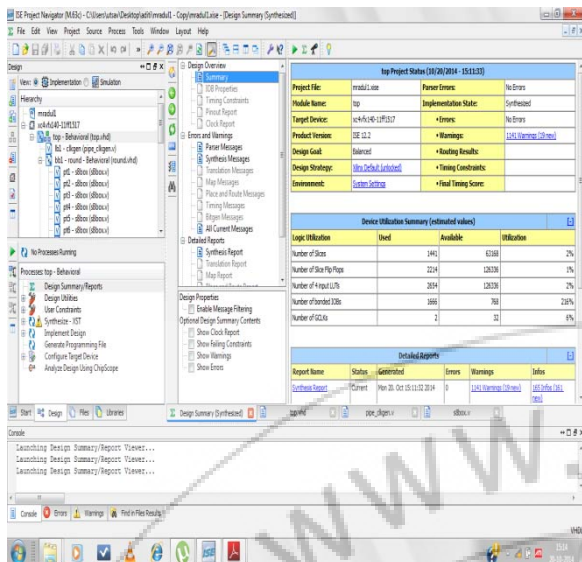


Figure 3 synthesize summary

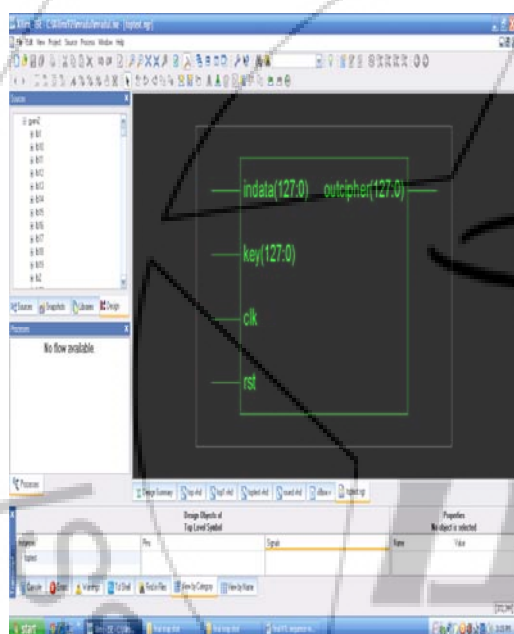


Figure 4 RTL view

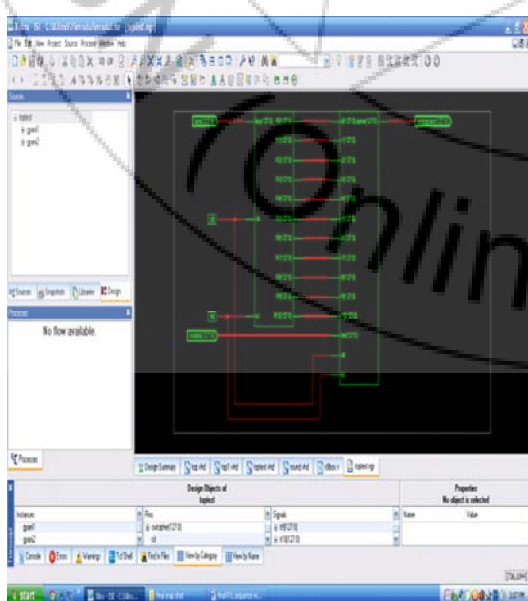


Figure 5: RTL view (Internal)

#### 4. Result Comparison

Table 2: Comparison of the implemented design of the module with reference design

Module	Our	Base paper 1	Base paper 2	Base paper 3	Base paper 4
Area ((Slices))	1441	3420	1853	4720	17450
Frequency (M Hz)	527.426	263.92	140.390	245	242.15
Throughput (MBPS)	1323	2815	352	2909	3090

#### 5. Conclusion

The proposed work has implemented several modules of advanced encryption standard and the simulation done on Xilinx ISE simulator. In the proposed work pipeline is introduce into the AES method. The target device FPGA vertex 4 has been used for the validation of the proposed AES.

In the proposed work number of slices used for advanced encryption standard is reduced and the throughput is increased, in this work we have used a pipeline for all 4 stages of AES, with the help of pipeline structure the area is reduced and time taken for encryption is also reduced. We used a register in between two stages then the area (slices) is reduced and throughput is increased compare to the reference design [1]. This AES is more secured compare to previous one because in this work the combination of key is more random.

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