Image Compression Approach for Encryption and Decryption using RGB Content with DCT Function

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Abstract: In this cryptosystem, we have considered RGB images for two-dimensional (2D) data security. Security of RGB images during transmission is a major concern, discussed globally. This paper proposes a novel technique for color image security by random cipher associated with 2D discrete cosine transform. Existing techniques have discussed the security of image data on the basis of the keys only (which provide only one layer of security for image data), but in the proposed cryptosystem, the keys and the arrangement of cipher parameters are imperative for correct decryption of color image data. Additionally, key multiplication side (pre or post) with. The RGB image data should inevitably be known, to correctly decrypt the encrypted image data. So, the proposed cryptosystem provides three layers of security for RGB image data. A computer simulation on standard examples and results is given to support the fixture of the scheme. Security analysis and detailed comparison between formerly developed techniques and proposed cryptosystem are also discussed for the robustness of the technique. This method will have large potential usage in the digital RGB image processing and the security of image data. There are many image encryption schemes have been proposed, each one of them has its advantages and disadvantages. This paper presents practical approach on image encryption and decryption technique using matrix transformation. The proposed scheme is useful for encryption of large amounts of data, such as digital images. First, we use discrete cosine transformation to get a blocked image. Second, a pair of keys is given by using matrix transformation. Third, the image is encrypted using private key in its transformation domain. Finally the receiver uses the public key to decrypt the encrypted messages. This technique satisfies the characters of convenient realization, less computation complexity and good security. The salient features of the proposed image encryption method are loss-less, symmetric private key encryption, a very large number of secret keys, and key-dependent pixel value replacement.

Keywords: Data security; cipher; data encryption; data decryption; cryptography; discrete cosine transform

1. Introduction

Image encryption and decryption are essential for securing images from various types of security attacks. In this paper, we have proposed a first approach for an RGB image encryption and decryption using two stage random matrices affine cipher associated with discrete cosine transformation. Earlier proposed schemes for encoding and decoding of images discussed only about the keys, but in our proposed approach, keys and the arrangement of cipher parameters are mandatory. We have also formulated a formula for all the possible range to choose keys for encrypting and decrypting an RGB image [1]. Computer simulation with a standard example and result is given to analyze the capability of the proposed approach. We have given security analysis and comparison between our proposed technique and others to support for robustness of the approach. This approach can be used for transmission of image data efficiently and securely. Quantitative evaluation methods have also been suggested [2]. However, due to the difficult nature of the problem, there are few automatic algorithms that can work well on a large variety of data. Security of image data in an insecure network is a major issue. Image data is highly sensitive and is prone to being decoded abruptly by intruders. Cryptographic systems are used extensively, to ensure secrecy and authenticity of sensitive information. Cryptography allows us to transmit data in such a way, that it is understood only at the receiver end. The original image data is called the plaintext data, which must be kept secure. The plaintext data changes into encrypted data by some algorithm[4]. This encrypted data is known as cipher text (encrypted image data), which is then transmitted through unsecured network. The procedure for secure transmission of RGB image data as well as the keys through insecure network is given in Fig. 1. At the receiver’s end, transmitted data is decrypted back into the plaintext. The aim of cryptography is to ensure high end communication between the senders and receiver without any loss of information. But cryptanalysts try to break the security of data, and this process is known as hacking. Security, refers to the following aspects-confidentiality, data integrity, authentication and non-repudiation. In current time, maintaining the fidelity, security, and confidentiality of original image data is a critical issue. The security of image data, designed by existing techniques completely depends on the encryption key. If the attacker knows about the exact encryption key of these techniques, then he/she can recover the original image from encrypted image data. So, in these approaches, the security of encryption keys is also required. The technique proposes novel cryptosystem for RGB images by cipher with discrete Cosine transform (DCT). In the approach [3] each component of the RGB image data is separated, and then the algorithm is applied to it. The scheme provides security of image data by the keys and the arrangement of RMAC parameters, which is two layer securities (first layer security by keys and second layer security by the arrangement of RMAC parameters). Various schemes have been developed for security of image data; such as: the techniques propose security of image data using Fourier transform; the schemes have given security of images in gyrator transform domain; the cryptosystems have developed image encryption using Hartley transform; the
techniques have presented image coding by wavelet transform; the approach have also proposed security of image data. According to, recent studies for the security of RGB image data, some attacks such as: brute-force attack, cropping attack, known-plaintext attack, chosen-plaintext attack, and chosen-cipher text attack, etc. can penetrate into the existing techniques for security (robustness) of image data [5].

2. Discrete Cosine Transformation & Approach

Research in image security was first motivated, in part, by the increasing use of digital means to transmit, store and view entertainment information such as images and video. The digital form allowed the perfect duplication of information and almost seamless manipulation and tampering of the data [10]. This created new types of security attacks not (as seriously) addressed in the past by the entertainment industry. The paradigm shift from analog to digital multimedia for entertainment has had an enormous impact for artists, publishers, copyright holders and consumers alike providing flexible and more accessible business models. In such a setting, one natural question that arises is the security and confidentiality of a digital packet of multimedia information.

To addresses the problem of reducing the memory space and amount of data required to represent a digital image. Image compression plays a crucial role in many important and adverse applications and including televideo conferencing, remote sensing, and document, medical and facsimile transmission. The need for an efficient technique for compression of Images ever increasing because the raw images need large amounts of disk space seems to be a big disadvantage during transmission & storage. Even though there are so many compression techniques already available-a better technique which is faster, memory efficient and simple surely suits the requirements of the user. In this paper the Spatial Redundancy method of image compression using a simple transform technique called Discrete Cosine Transform is proposed. This technique is simple in implementation and utilizes less memory. A software algorithm has been developed and implemented to compress the given RGB image using Discrete Cosine Transform techniques in a MATLAB platform.

![Figure 1: Transmission of secure data though an insecure network](image1)

![Figure 2: Analytical Approach](image2)

![Figure 3: Image encryption scenario](image3)

![Figure 4: Division of Source Image into Planes](image4)

The image preparation is NOT BASED on
1. 9-bit YUV encoding
2. Fixed number of lines and columns
3. Mapping of encoded chrominance

Source image consists of components (C) and to each component we assign YUV, RGB or TIQ signals."
d. Making the inverse DCT transformation and uniting all \( P \times P \) blocks, the final result is defined as

\[
X_2(M \times N) = IDCT \{X(M \times N)\} 
\]

(4)

e. Keeping all the transformed coefficients between 0 and 1 [5].

i. Get the minimum of \( X_2(M \times N) \) named as \( \text{Min} \)

\[
\text{Min} = \max((-1) \times X_2(M \times N)) 
\]

ii. Ensure all the coefficient of \( X_2(M \times N) \) more than 0

\[
\text{I.e. } X_2(M \times N) = X_2(M \times N) + \text{Min} 
\]

iii. Get the maximum of updated \( X_2(M \times N) \) named as \( \text{Max} \)

\[
\text{I.e. } \text{Max} = \max(X_2(M \times N)) 
\]

iv. Ensure all the coefficient of \( X_2(M \times N) \) less than 1

\[
\text{I.e. } X_2(M \times N) = X_2(M \times N) / \text{Max} 
\]

\[\text{Max} = x_1 \times \text{Max} \]

2.2 Image Preparation - Blocks

Images are divided into data units, called blocks – definition comes from DCT transformation since DCT operates on blocks. Lossy mode – blocks of 8x8 pixels; lossless mode – data unit 1 pixel

Decryption

The decryption operation is a usual correlation process with five elements: (1) block length \( P \) (2) encryption matrix dimension \( K \) (3) public key \( A-tU \) (4) the coefficient minimum \( \text{Min} \) (5) the coefficient maximum \( \text{Max} \). Suppose \( X3M\times N \) denotes the encrypted image, the details of decryption are following [6]:

- Recovering all coefficients of \( X3 \) M\( \times N \)

\[
X3 M \times N = X3 M \times N \times \text{Max} - \text{Min} 
\]

(6)

- Applying DCT transformation to each distinct \( P \times P \) block of \( X3 \) M\( \times N \)

\[
X4 M \times N = \text{DCT}(X3 M \times N, [P P]) 
\]

(7)

- Decrypting the frontal \( P \times K \) coefficients of every \( P \times P \) block, respectively.

Let \( D1 \) denotes the matrix composed by \( P \times K \) coefficients of certain \( P \times P \) block \( D0 \), the corresponding decryption data \( D2 \) \( 2R\times K \) by using the public key \( A-t U \) can be given as following:

\[
\text{D} 2 = (A-t U) t D1 \\
\text{D} 2 = (U-t A^{-1}) (A U) X0 \\
\text{D} 2 = (U-t U) X0 
\]

(8)

Because the column vector of \( U \) is a set of orthonormality bases, it is easily proved: \( U t U = E \). So, we can draw the conclusion:

\[
\text{D} 2 = X0 
\]

(9)

- Replacing the frontal \( P \times K \) coefficients of \( D0 \) with \( D2 \) and 0 ...

(10)

- Making the inverse DCT transformation and uniting all \( P \times P \) blocks, the final result is defined as \( X5M\times N \).

\[
X5 M \times N = \text{IDCT}(X4 M \times N) 
\]

(11)

- Saving the decrypted image as jpg file.

2.3 Data Unit Ordering

- Non-interleaved: scan from left to right, top to bottom for each color component
- Interleaved: compute one “unit” from each color component, then repeat full color pixels after each step of decoding but components may have different resolution

\[\text{Set your page as A4, width 210, height 297 and margins as follows [3]:} \]

\[
\text{Left Margin 17.8 mm (0.67")} \\
\text{Right Margin 14.3 mm (0.56)} \\
\text{Top Margin – 17.8 mm (0.7")} \\
\text{Bottom Margin – 17.8 mm (0.7")} 
\]

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You should use Times Roman of size 10 for all fonts in the paper. Format the page as two columns:
- Column Width 86.8 mm (3.42”)
- Column Height – 271.4 mm (10.69”)
- Space/Gap between Columns - 5.0 mm (0.2”).

3. Image encryption- Decryption Process

Shift values \([0, 2^p - 1]\) to \([-2^{p-1}, 2^{p-1} - 1]\)  
  
- e.g. if \((P=8)\), shift \([0, 255]\) to \([-127, 127]\)  
- DCT requires range be centered around 0

Values in 8x8 pixel blocks are spatial values and there are 64 samples values in each block:
- Forward DCT Convert from spatial to frequency domain.
- convert intensity function into weighted sum of periodic basis (cosine) functions  
- identify bands of spectral information that can be thrown away without loss of quality  
- Intensity values in each color plane often change slowly.

In the proposed method, a comparative study of selective image encryption using DCT with Stream Cipher is done. In the DCT method, the basic idea is to decompose the image into 8x8 blocks and these blocks are transformed from the spatial domain to the frequency domain by the DCT. Then, the DCT coefficients correlated to the lower frequencies of the image block are encrypted using the RC4 Stream Cipher. The concept behind encrypting only some selective DCT coefficients (the coefficients \([0,0], [0,1], [0,2], [1,0], [2,0], [1,1]\)) is based on the fact that the image details are situated in the lower frequencies and the human is most sensitive to the lower frequencies than to the higher frequencies. An extra security has been provided to the resulted encrypted blocks by shuffling the resulted blocks using the Shuffling Algorithm. Fig. 1.5 shows the general block diagram of the proposed method of selective image encryption.

3.1. Algorithm to Encrypt Image

Input: Target Image to be encrypted and the stream Key values.  
Output: Encrypted Image

Begin
- Read the image header, save the height of the image in variable height & the width in variable width and save the body image in an array image body.  
- Obtain how many blocks exist in an image row and how many ones in the column, by dividing the width and height of the image by \(N\), where \(N\) is equal to 8 (the required block size).

No Row B = Image Height / \(N\);  
No Col B = Image Width / \(N\);  
For all blocks in the image perform the following:
- Get_block (row_no, col_no)  
- Perform a DCT on the block and save the resulted coefficients in an array.  
- Round the selected coefficients, convert the selected coefficients to 11 bits; the 12th bit is used to save the sign of the coefficient.  
- Encrypt the selected coefficients by XORing the generated bit stream from the RC4 + Key with the coefficient bits, the sign bit of the selected coefficients will not be encrypted.  
- Perform an Inverse Discrete Cosine Transform (IDCT) and get the new block values and the resulted values could be positive or negative values due to the encryption step.  
- Apply the proposed shuffling algorithm on the resulted blocks to obtain the encrypted image.

End

3.2. Algorithm to Decrypt Image

Input: Target Image to be decrypted and the Encryption Key
Output: Original Image

Begin  
- Read the image header, save the height of the image in variable height & the width in variable width and save the body image in an array image body.  
- Obtain how many blocks exist in an image row and how many ones in the column, by dividing the width and height of the image by \(N\), where \(N\) is equal to 8 (the required block size).

No Row B = Image Height / \(N\);  
No Col B = Image Width / \(N\);  
For all blocks in the image perform the following:  
- Get_block (row_no, col_no)  
- Perform a DCT on the block and save the resulted values in an array.  
- Round the selected coefficients, convert the selected coefficients to 11 bits; the 12th bit is used to save the sign of the coefficient.  
- Decrypt the resulted bits by using the generated bit stream from the RC4 + Key, by performing an XOR operation, the sign bit of the selected coefficients will remain.  
- Convert the resulted bits into integer values, and join the sign (from the step above) with each integer, if the coefficient is negative multiply it by \(-1\).  
- Perform an IDCT and get the new blocks.  
- Reshuffle the block, since the shuffling algorithm generates the same row and column numbers to return the shuffled blocks into their original locations.  
- Reconstruct the image to get the original Image.

End

3.3. Shuffling Algorithm

Input: Key, number of blocks in the row \((\text{NoRows})\), number of blocks in the column \((\text{NoCols})\) and the resulted encrypted image saved in an array.

Output: A new shuffled image

Begin  
For \(i = 0\) to \((\text{NoRows} \times \text{NoCols})\)  
- NewVal \([i]\) = \((\text{Key} \times i) \mod (\text{NoRows} \times \text{NoCols})\)  
End

For  
- \(k = 0\)  
- \(i = 0\) to \((\text{NoRows} \times \text{NoCols})\)  
- MoveBlock \((\text{ImageBlk}[i]), \text{ImageBlk}[k])\)  
- \(k++\)
4. Result and Outcomes

![Figure 9: Compression window with RGB Content](image9)

![Figure 10: DCT compression with color image](image10)

![Figure 11: Compression of image with Black & White image and RGB Content Image](image11)

5. Conclusion and Future Scope

In this work we proposed a new image encryption and compression method based on Embedding and Discrete Cosine Transform (DCT) using RGB content. For encryption, DCT blocks of transmitted images are rotated and mixed with a random image to hide them.

In the decryption stage, the covered images can be extracted from the mixtures by applying extraction algorithm. Finally using rotation keys and inverse discrete cosine transform, the original images can be reconstructed.

Therefore we can achieve a fast and secure image transmission. As a result of several computer simulations, the behavior of the proposed approach is confirmed. In this paper color images e used as original images, but grey color images can be applied in the same way.

Our future works include a more secure encryption method with an alternative rotation method and a reconstruction key. More complex rotation manner makes it harder for unauthorized people to reconstruct images without keys.

References

[1] Anita khandelwal and Bhavneesh Malik,” Transmission and Reception of image using encryption and decryption technique” international Journals of research in
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